

ECE 4160/5160

MAE 4910/5910

Prof. Kirstin Hagelskjær Petersen

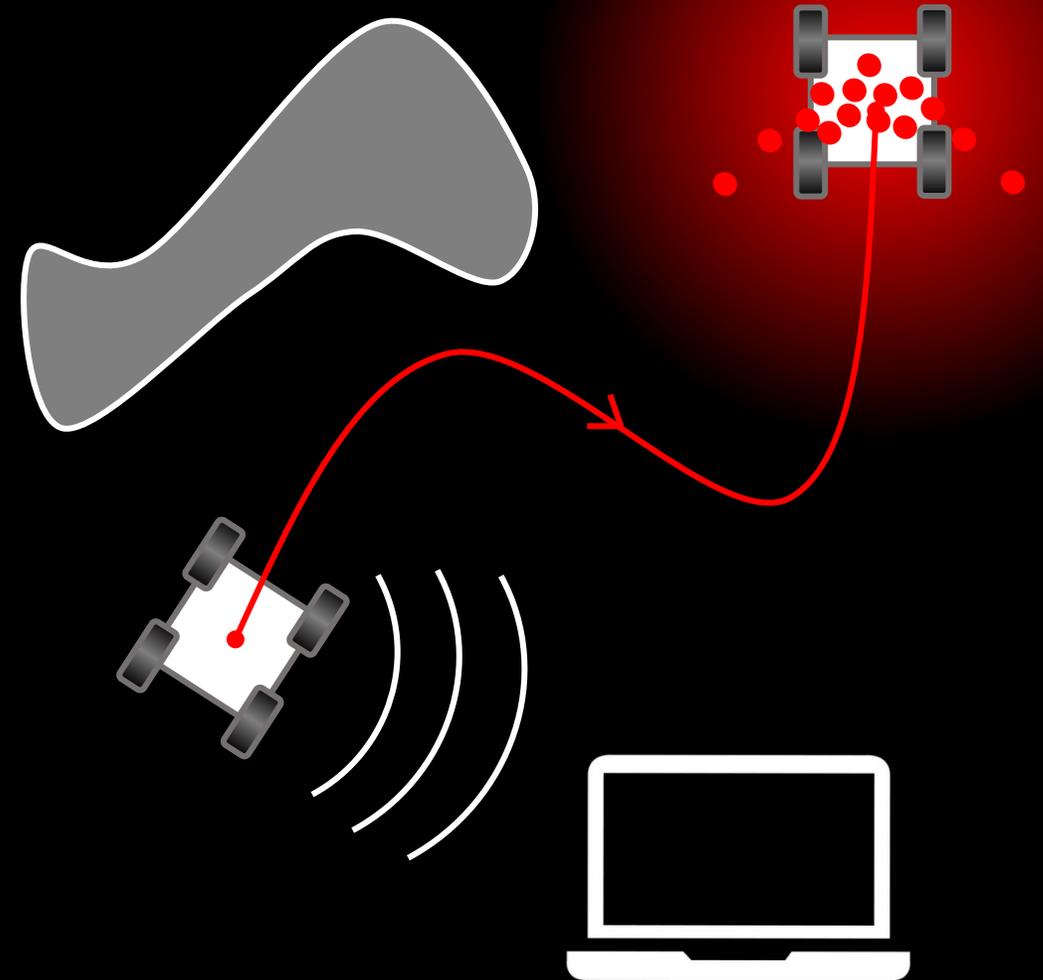
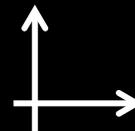
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Map Representations, Graphs and Graph Search

Outline of the next module on Navigation

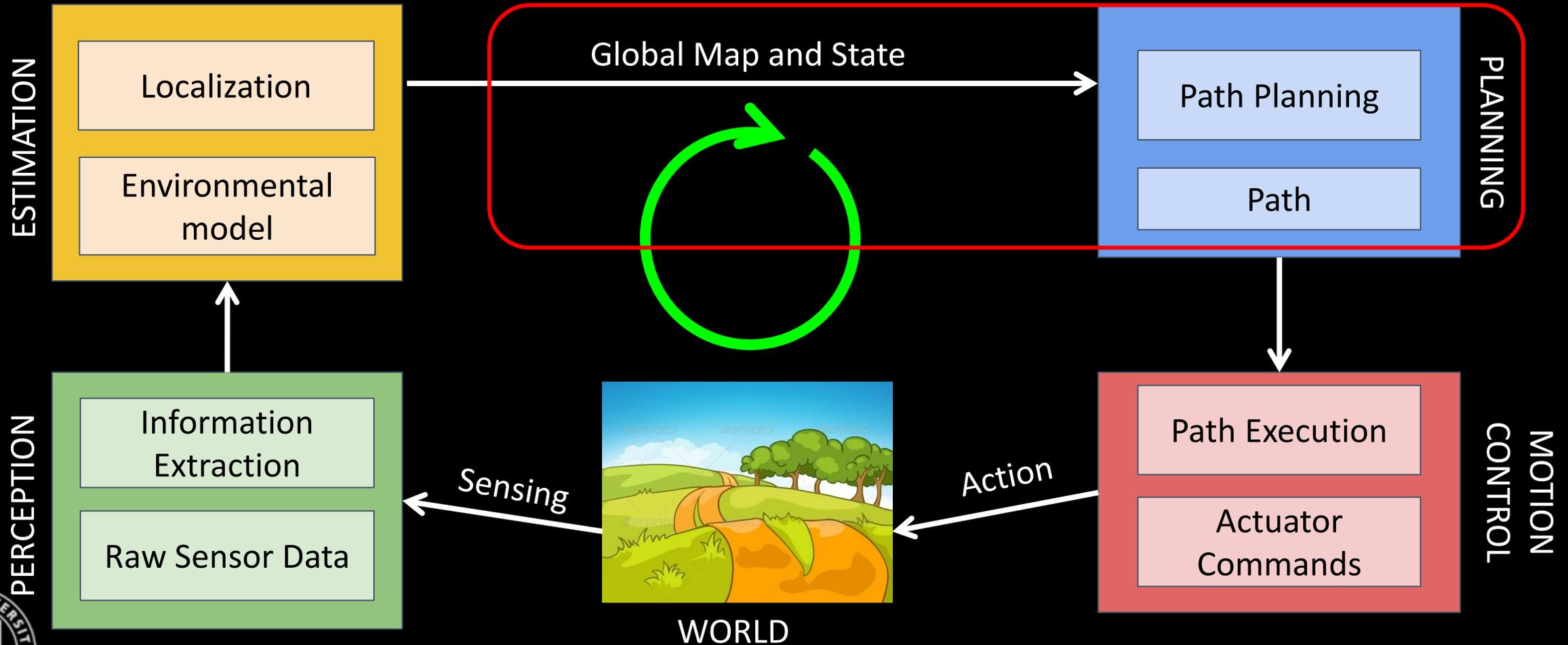
- Local planners
- Global localization and planning
 - Map representations
 - Continuous
 - Discrete
 - Topological
 - Maps as graphs
 - Graph Search Algorithms
 - Breadth First Search
 - Depth First Search
 - Dijkstras
 - A*

- Localization, Sensor and motion models, SLAM



Navigation

- Navigation breaks down to: Localization, Map Building, Path Planning

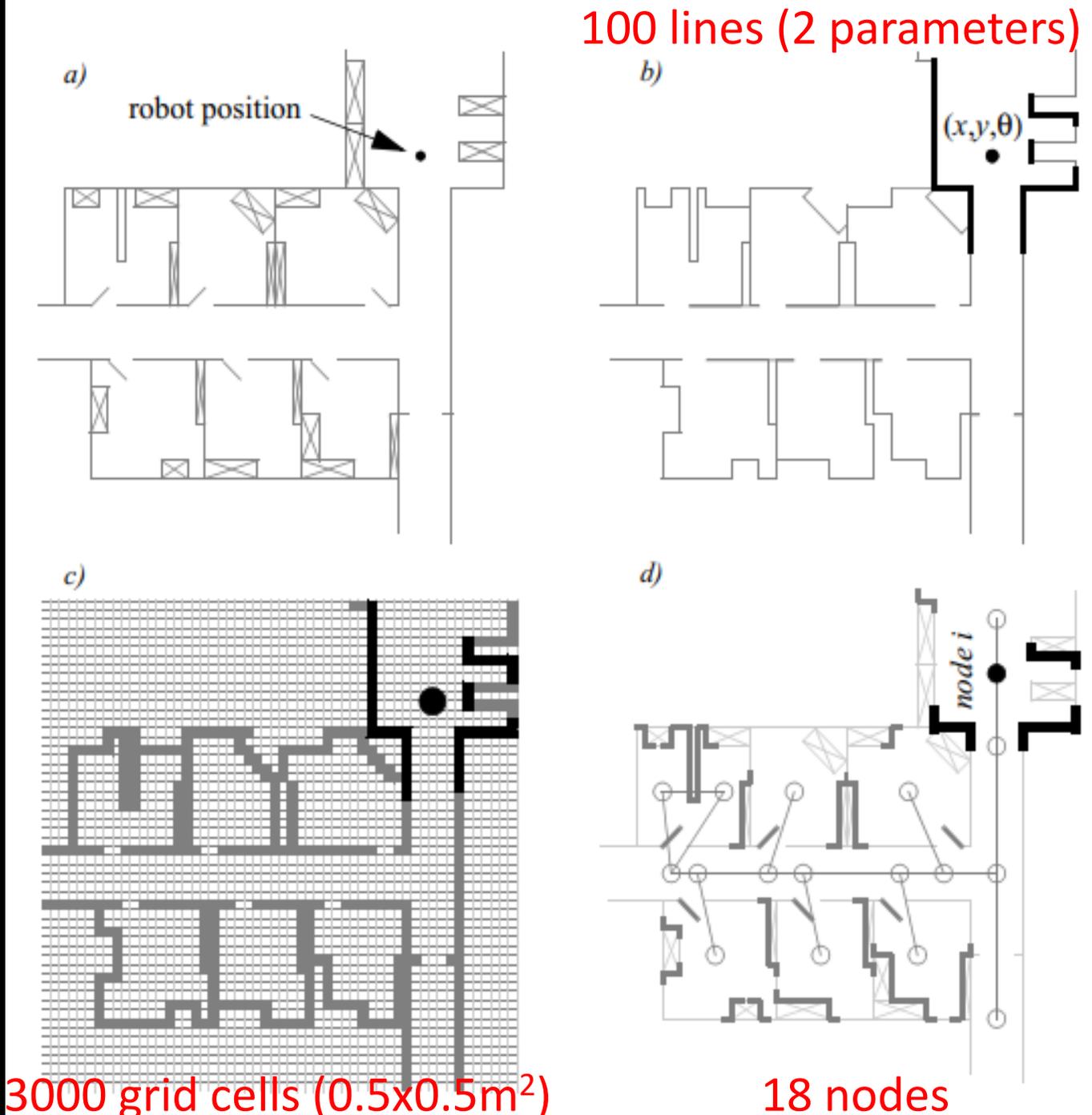


Map Representation

- (a) Building plan
- (b) line-based map
- (c) occupancy grid-based map
- (d) topological map

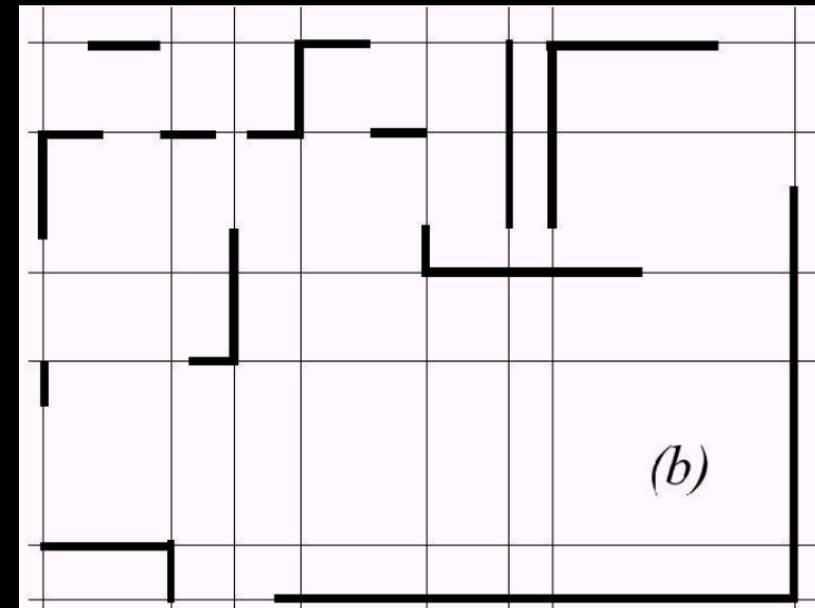
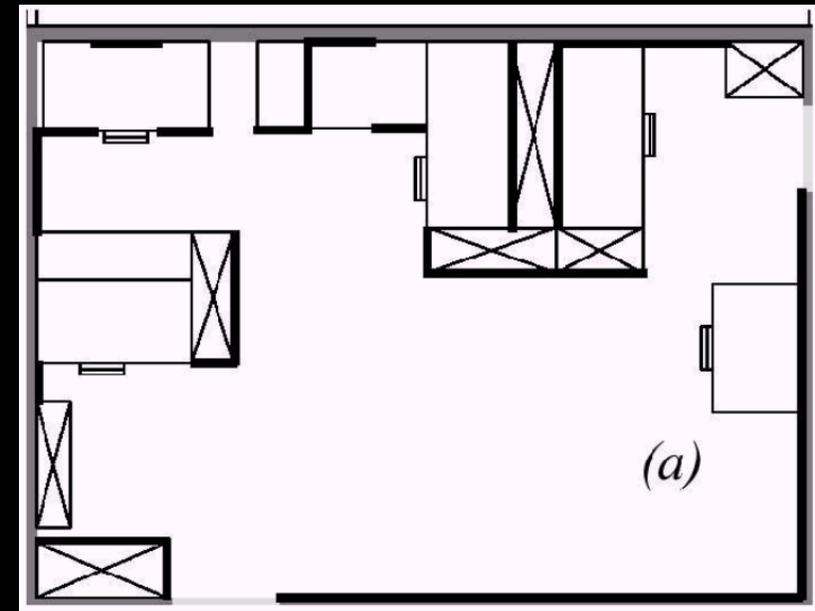
Important properties

- Memory allocation
- Computation



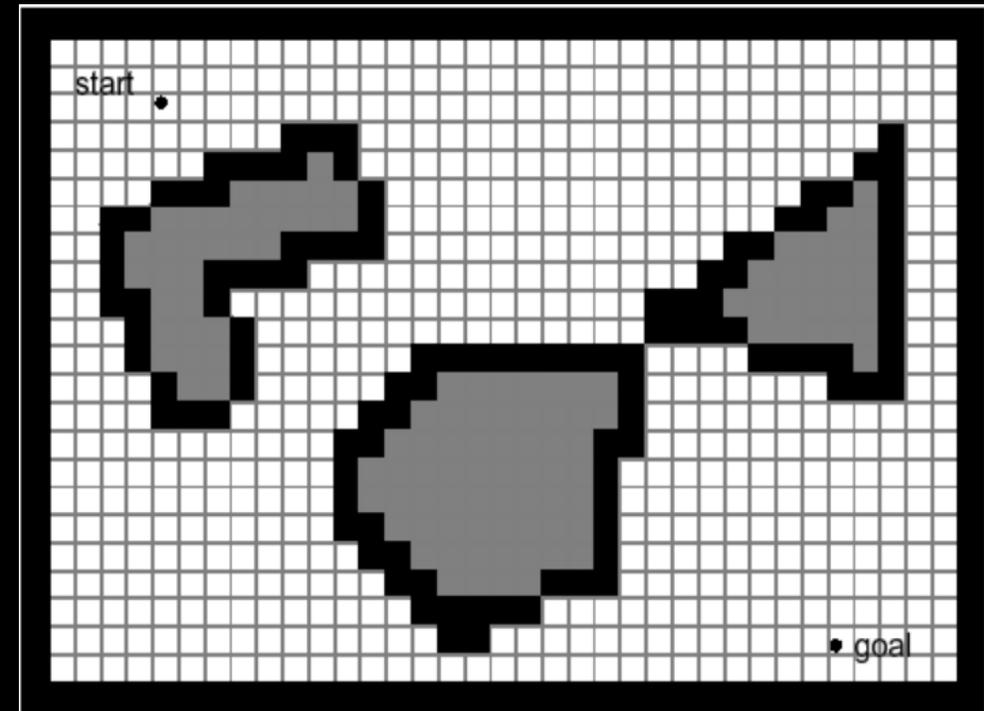
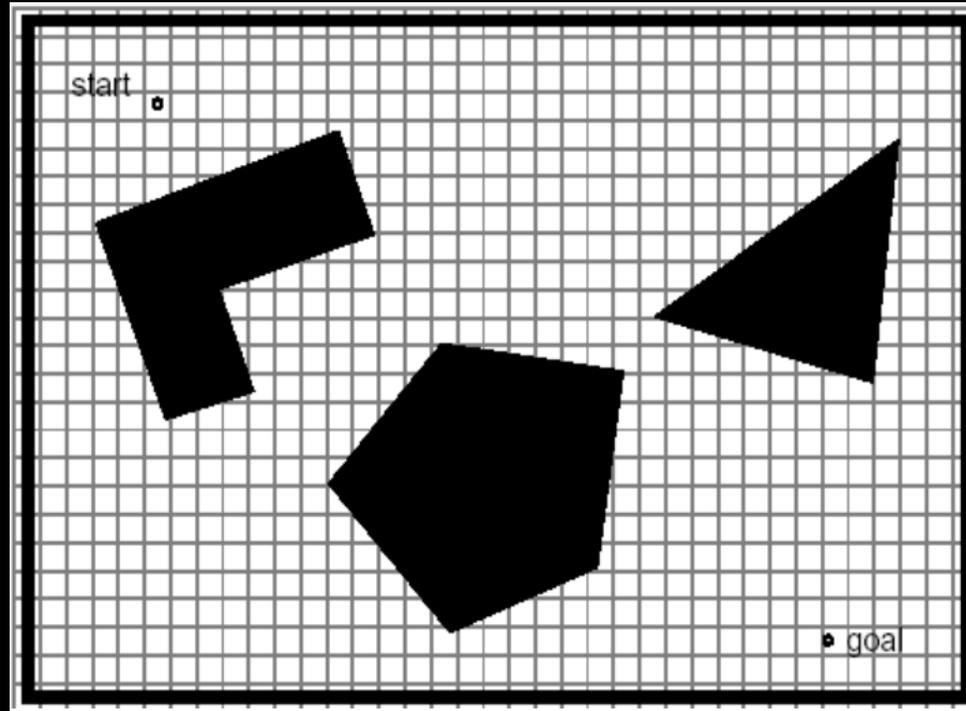
Continuous Representations

- Exact decomposition of the environment
- Used mainly in 2D representations
- Closed-world assumption
- Storage proportional to object density
- Example: Continuous line representations
 - Using range finders, we can extract lines/line segments in the environment

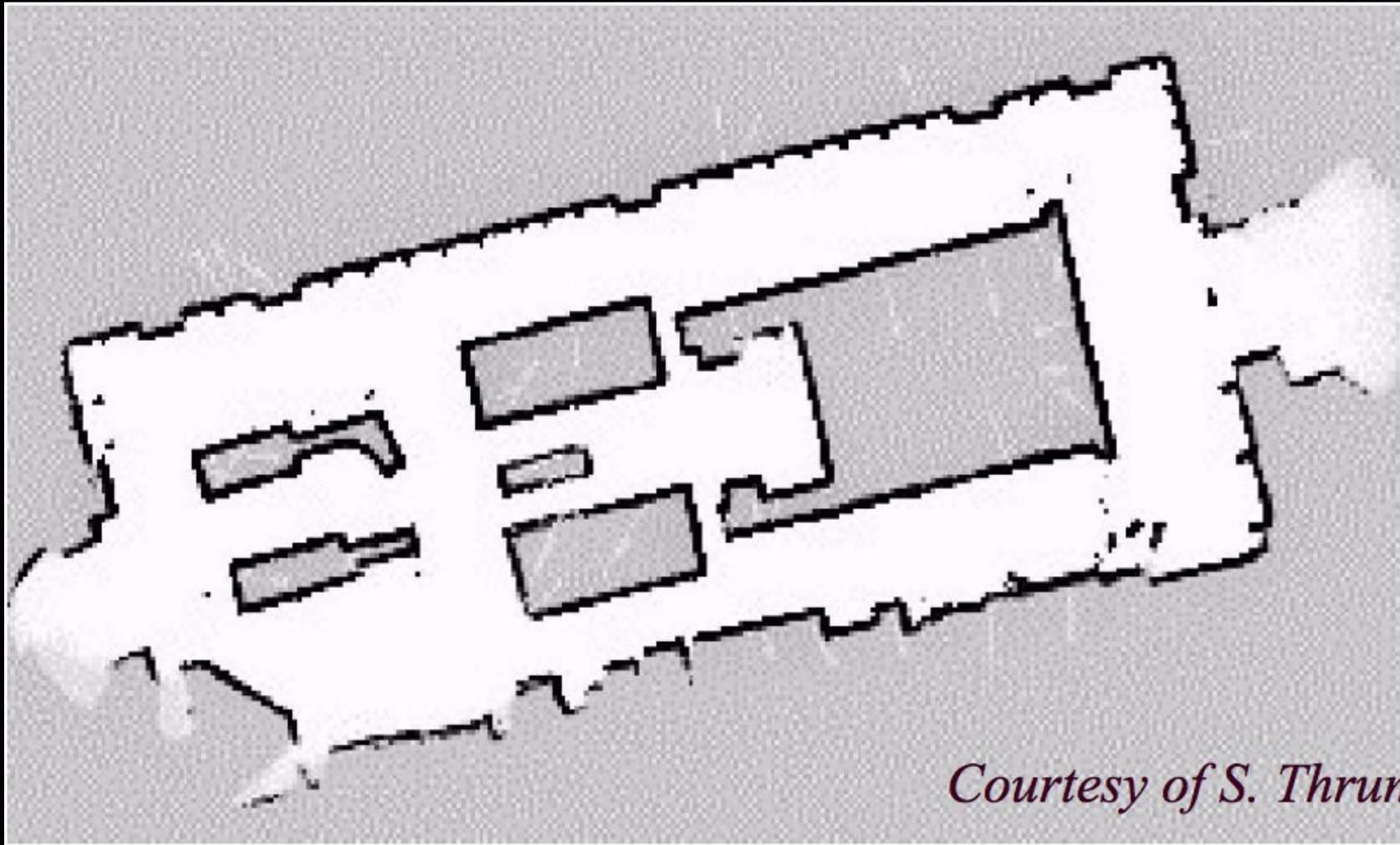


Fixed Decomposition

- Tessellate the world at a fixed resolution
- Approximate features given the resolution
- Most commonly used: Occupancy grid



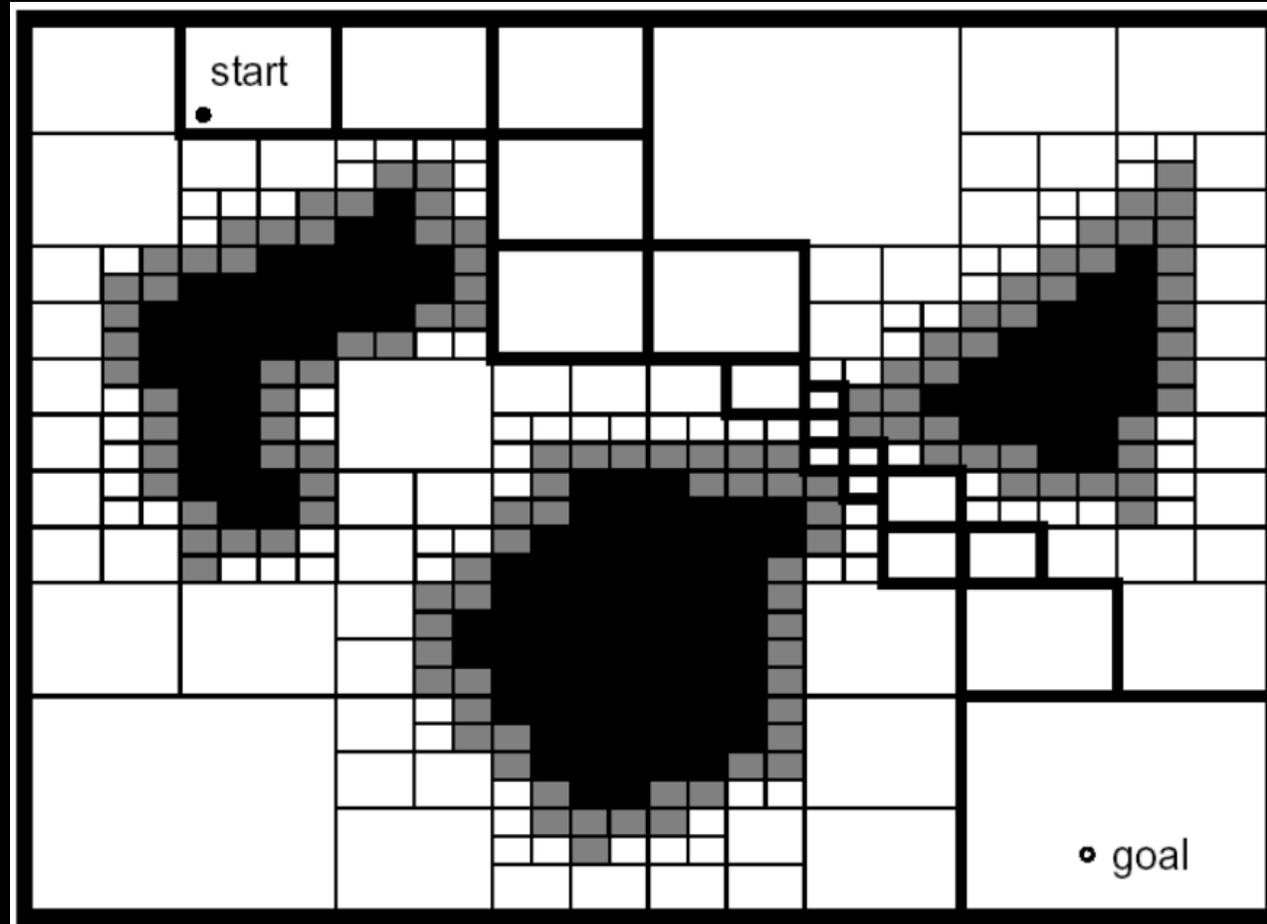
Fixed Decomposition



Courtesy of S. Thrun

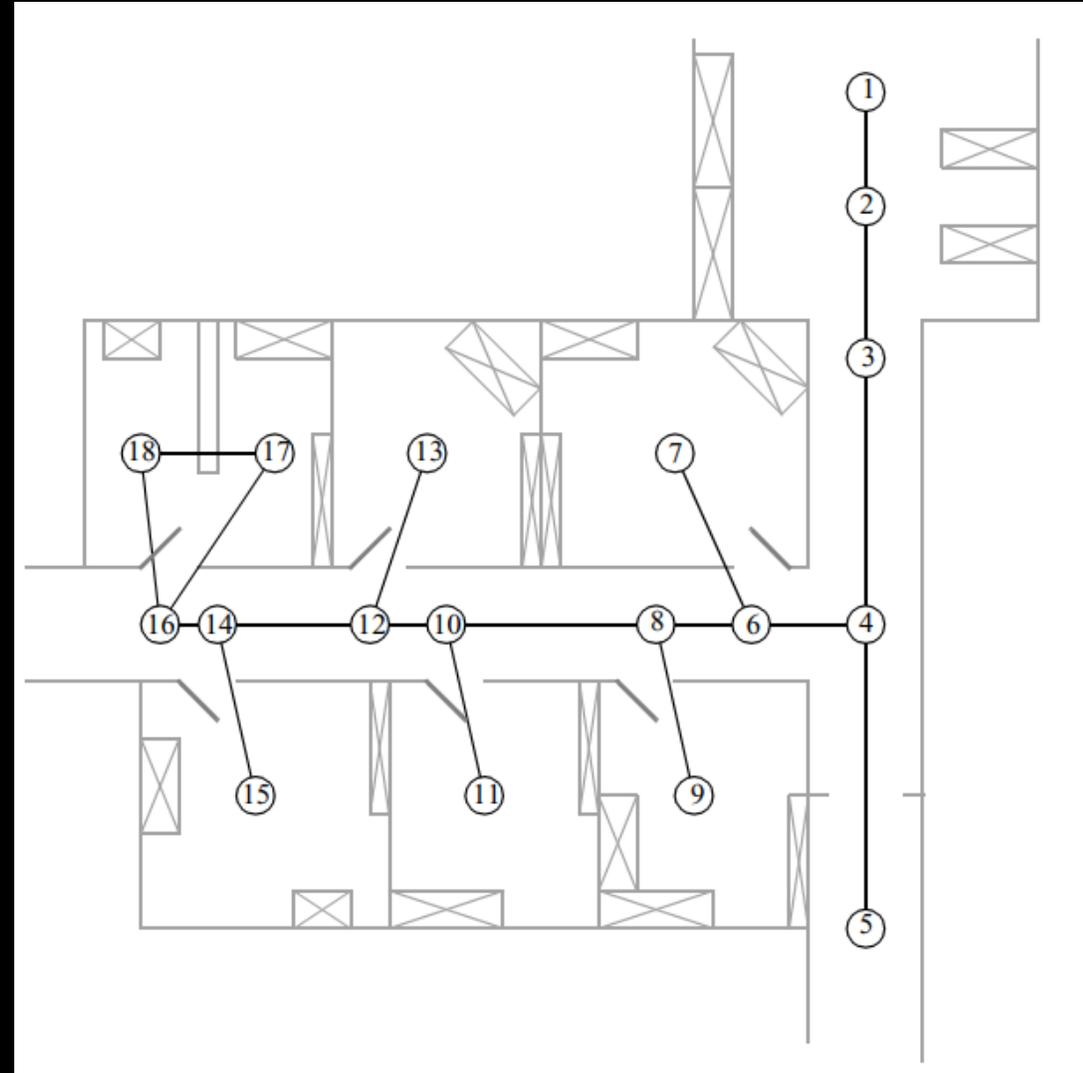
Adaptive Cell Decomposition

- Adapt cell size to features



Topological Decomposition

- A topological representation is a graph that specifies nodes and edges
 - Nodes denote areas in the environment
 - Edges describe environment connectivity
- Robots can...
 - ...detect their current position in terms of the nodes of the topological graph
 - ...travel between nodes using robot motion



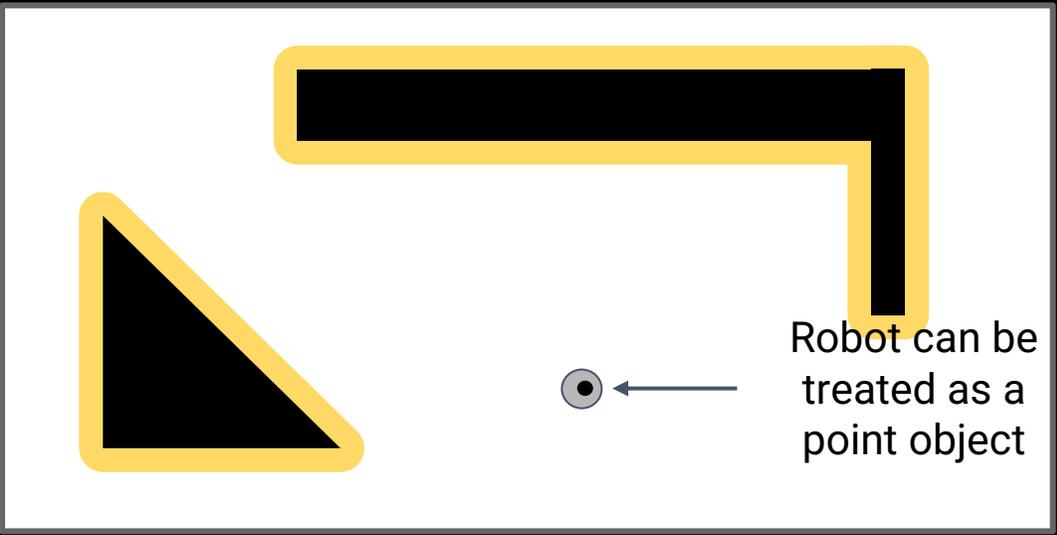
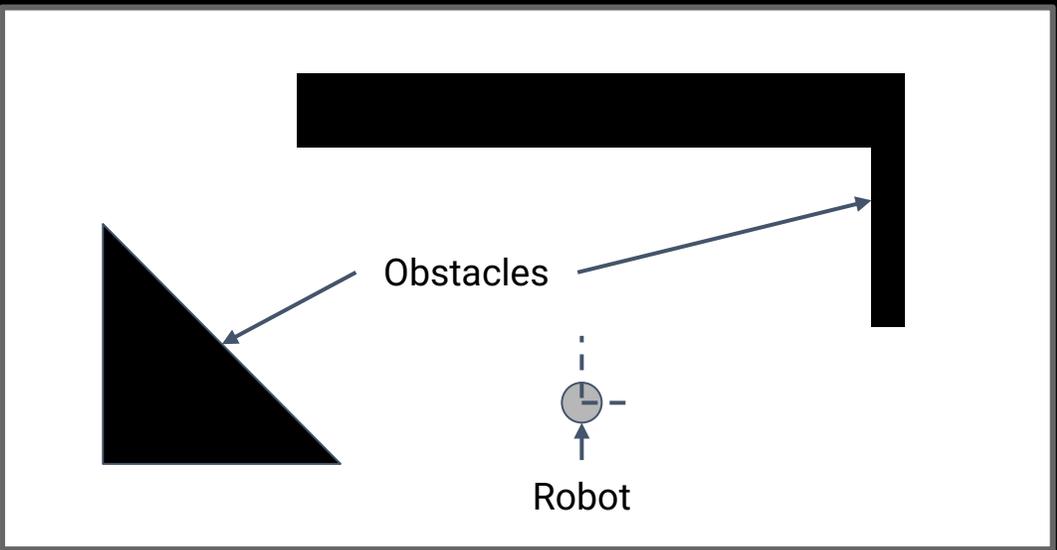
Topological Decomposition

- A topological representation is a graph that specifies nodes and edges
 - Nodes denote areas in the environment
 - Edges describe environment connectivity
- Robots can...
 - ...detect their current position in terms of the nodes of the topological graph
 - ...travel between nodes using robot motion
- Typical for 3D maps



How to represent the robot pose?

- Physical robots take up space
- Expand obstacles
- Represent maps in configuration space instead of Euclidean space



Map Representation Considerations

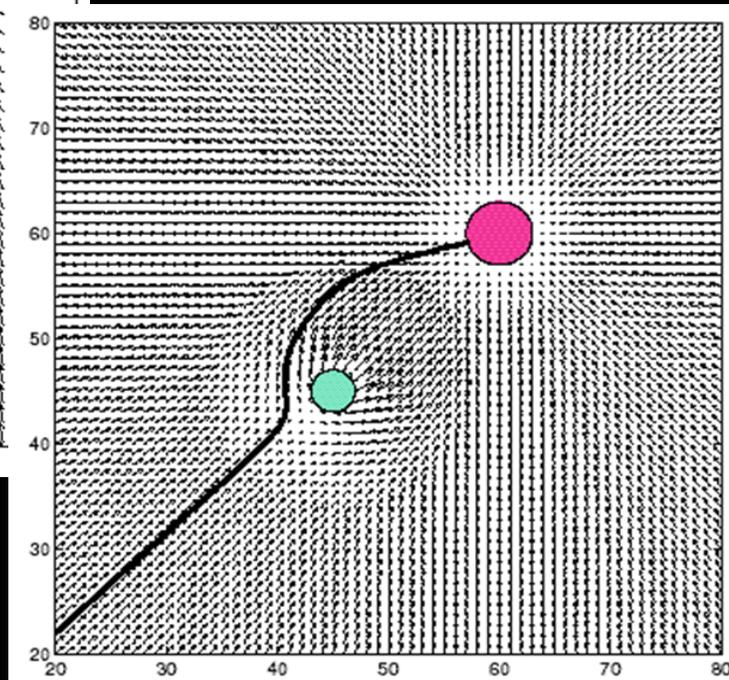
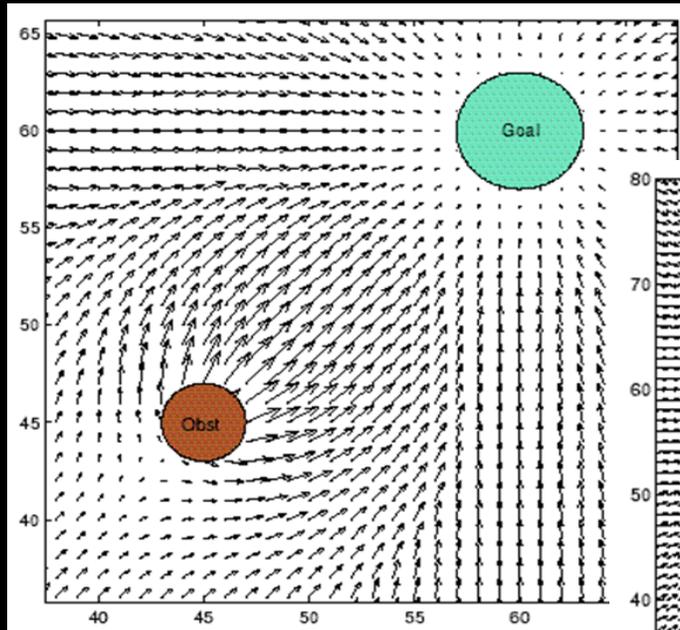
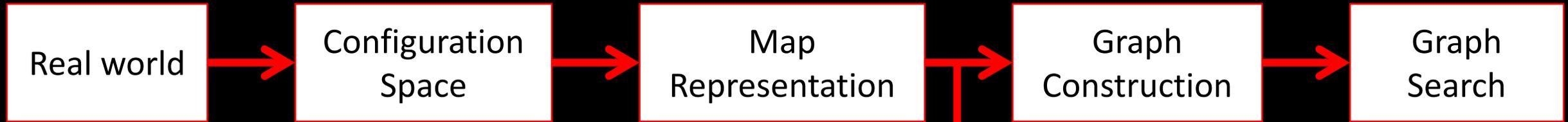
- The precision of the map must appropriately match the precision with which the robot needs to achieve its goals
- The precision of the map and the type of features represented must match the precision and data types returned by the robot's sensors
- The complexity of the map representation has direct impact on the computational complexity of reasoning about mapping, localization, and navigation

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Constructing Graphs

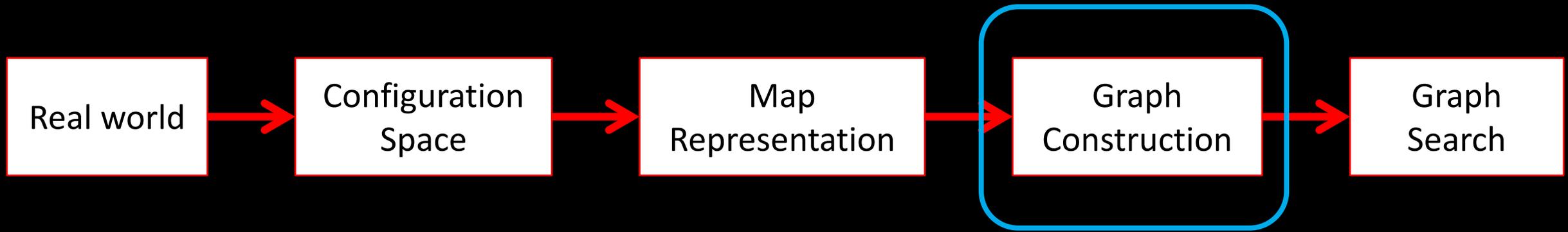
Modelling path planning as a graph search problem



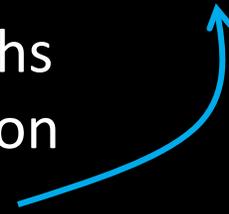
Common alternatives

- Optimal control
- Potential fields

Modelling path planning as a graph search problem



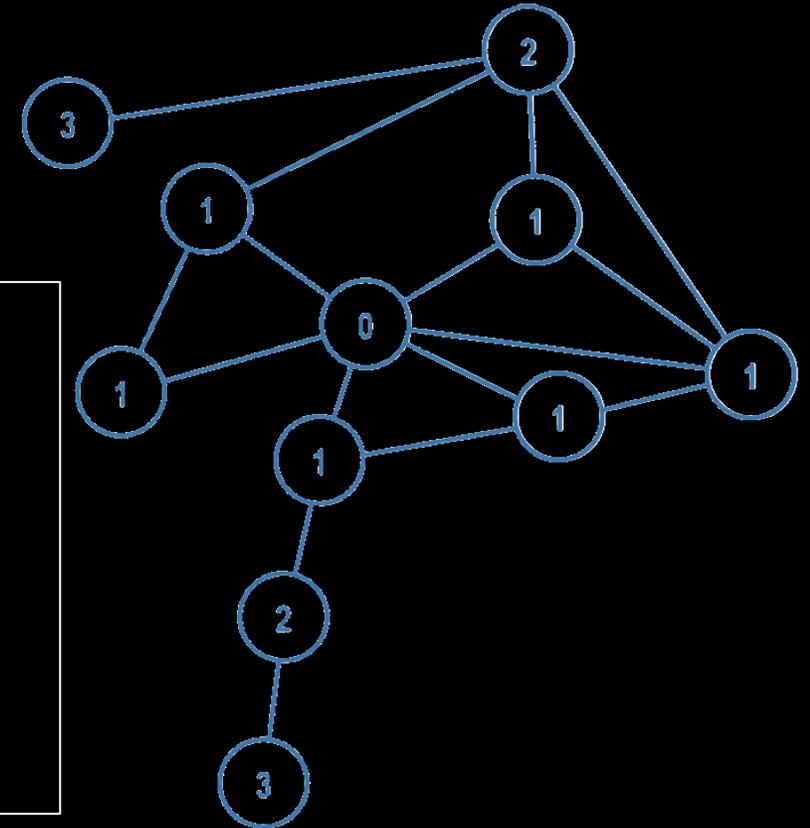
- Topological Graphs
- Cell decomposition
- Visibility Graphs
- RRT
- PRM



Graph Construction

- Transform continuous/discrete/topological maps to a discrete graph
- Why?
 - Model the path planning problem as a search problem
 - Graph theory has lots of tools
 - Real-time capable algorithms
 - Can accommodate for evolving maps

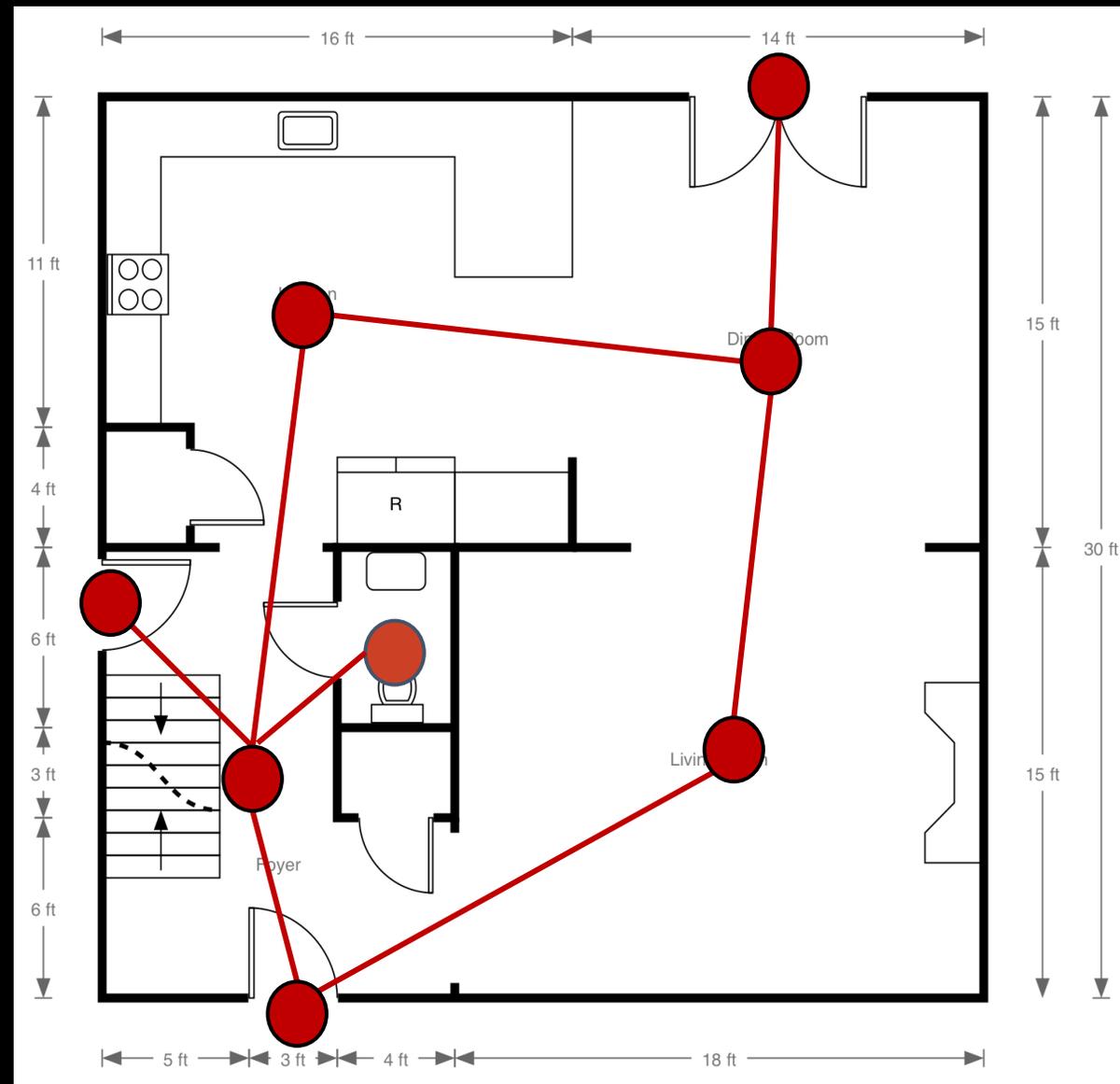
1. Divide space into simple, connected regions, or “cells”
2. Determine adjacency of open cells
3. Construct a connectivity graph
4. Find cells with initial and goal configuration
5. Search for a path in the connectivity graph to join them
6. From the sequence of cells, compute a path within each cell
 - e.g. passing through the midpoints of cell boundaries or by sequence of wall following movements



Geometry-Based Planners

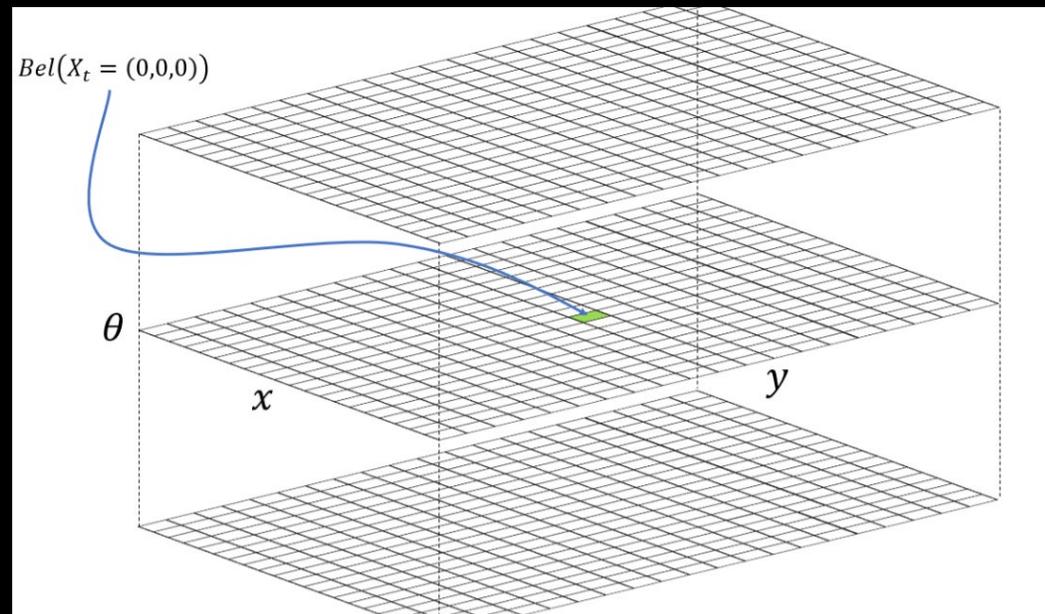
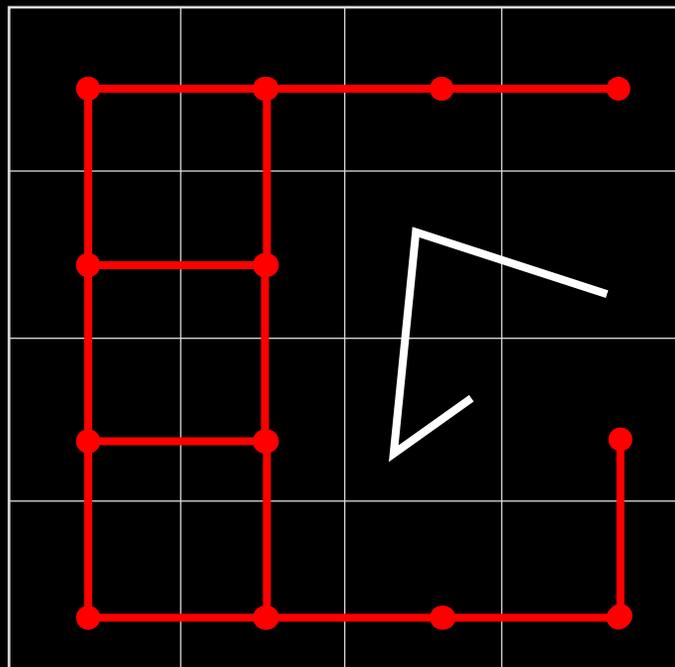
Topological Maps

- Good abstract representation
- Tradeoff in # of nodes
 - Complexity vs. accuracy
 - Efficient in large, sparse environments
 - Loss in geometric precision
- Edges can carry weights
- Con: Limited information

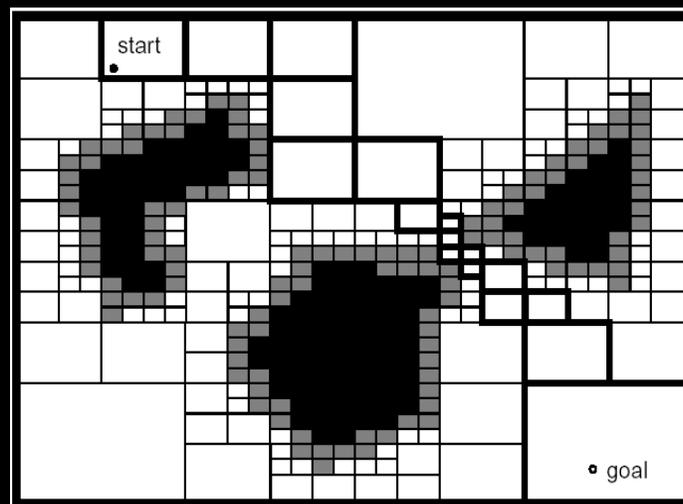


Fixed Cell Decomposition

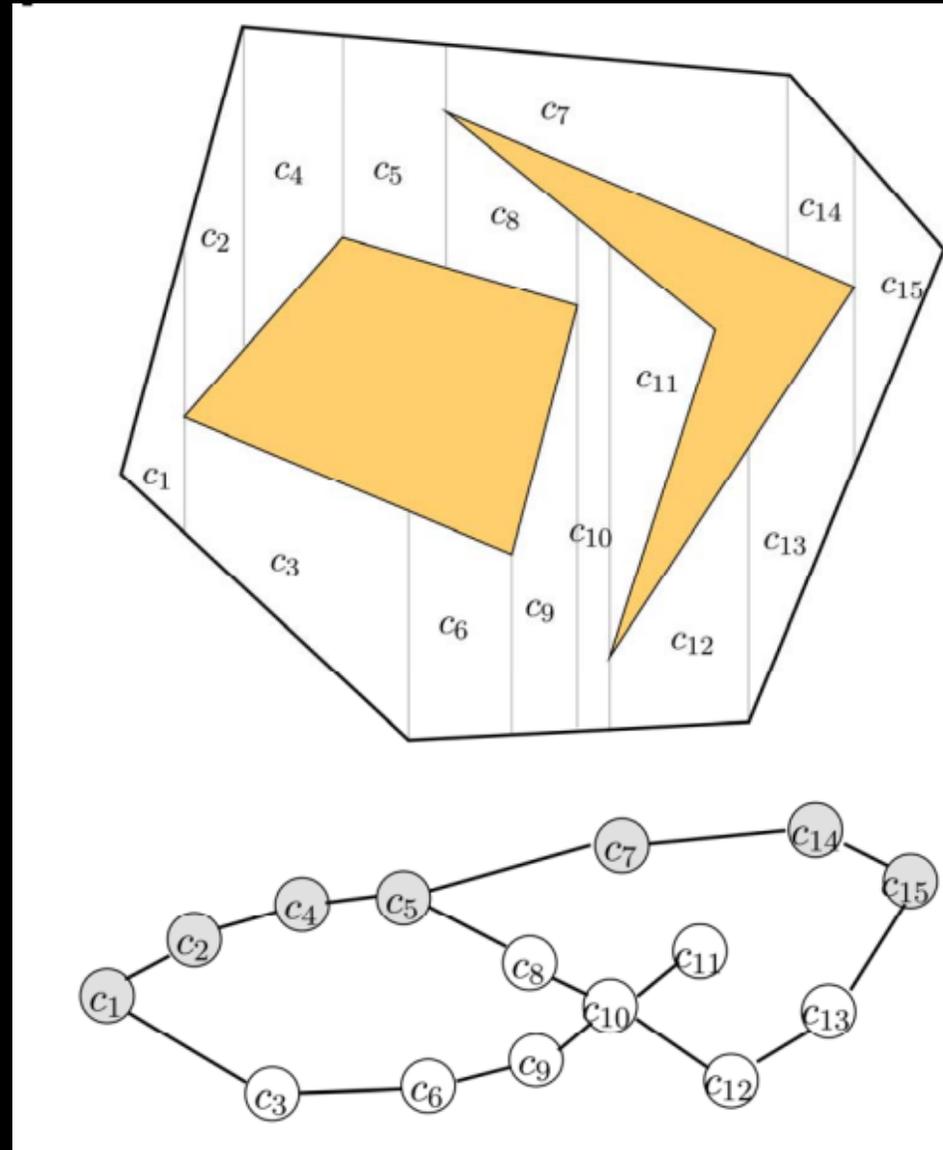
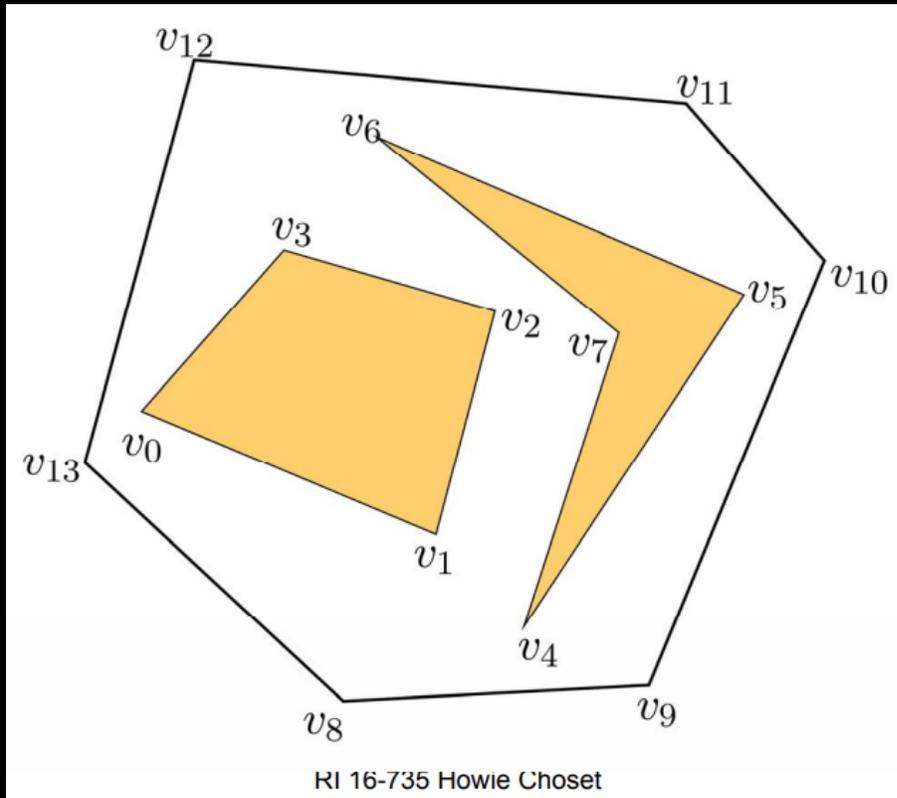
(Lab 9-12)



Adaptive Cell Decomposition

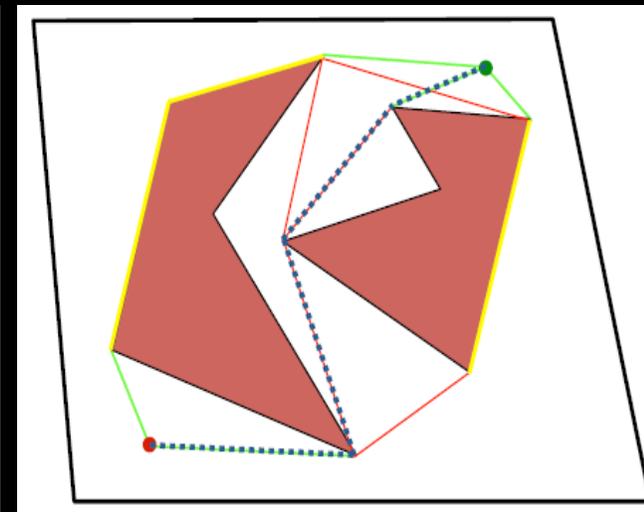
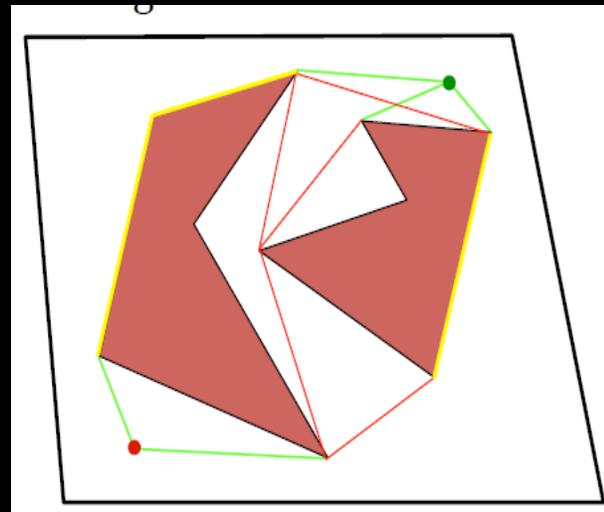
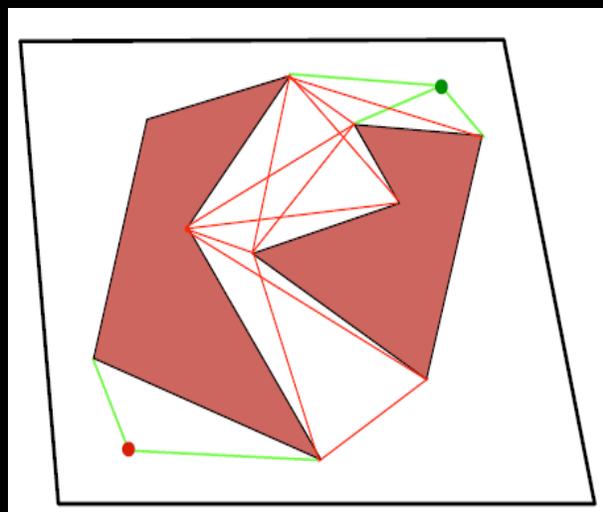
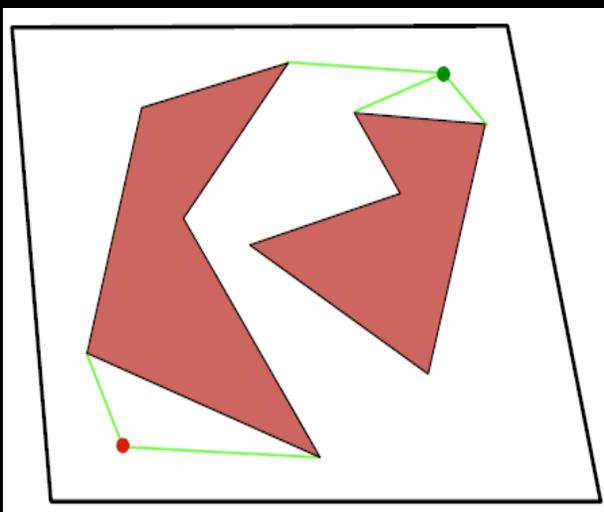


Trapezoidal Cell Decomposition



Visibility Graphs

- Connect initial and goal locations with all visible vertices
- Connect each obstacle vertex to every visible obstacle vertex
- Remove edges that intersect the interior of an obstacle
- Plan on the resulting graph

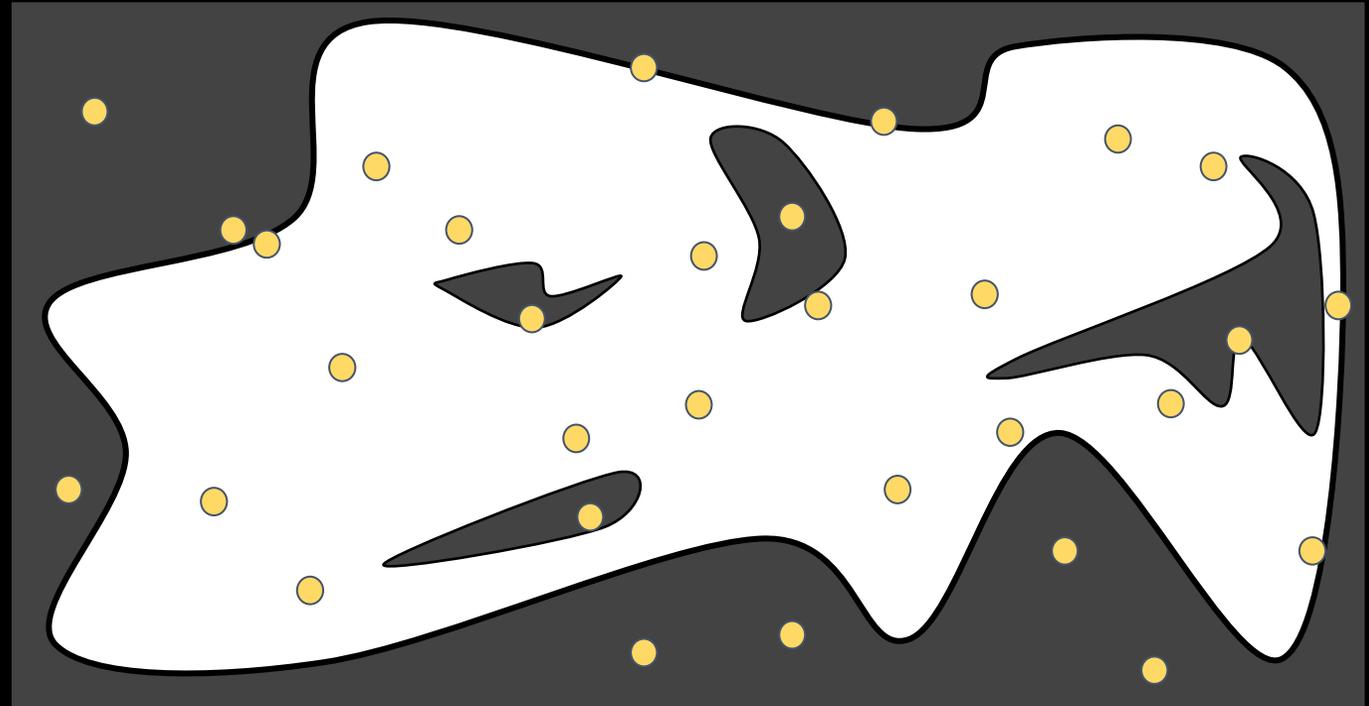


Sampling-Based Planners

- Rather than computing the C-Space explicitly, we sample it
- Often efficient in high dimensional spaces
- Compute if a robot configuration has collisions
 - Just requires forward kinematics
 - (Local path plans between configurations)
- Examples
 - Probabilistic Roadmaps (PRM)
 - Rapidly Exploring Random Trees (RRT)

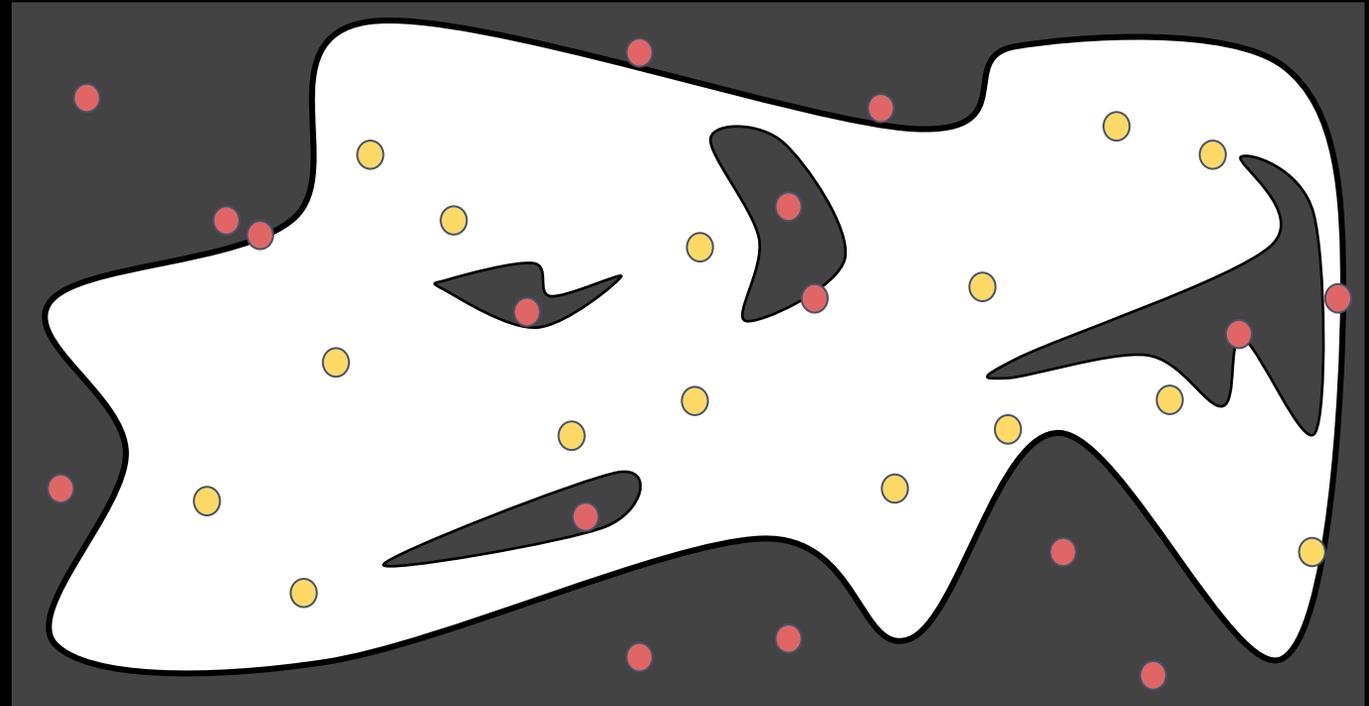
Probabilistic Roadmaps

- Configurations are sampled by picking coordinates at random



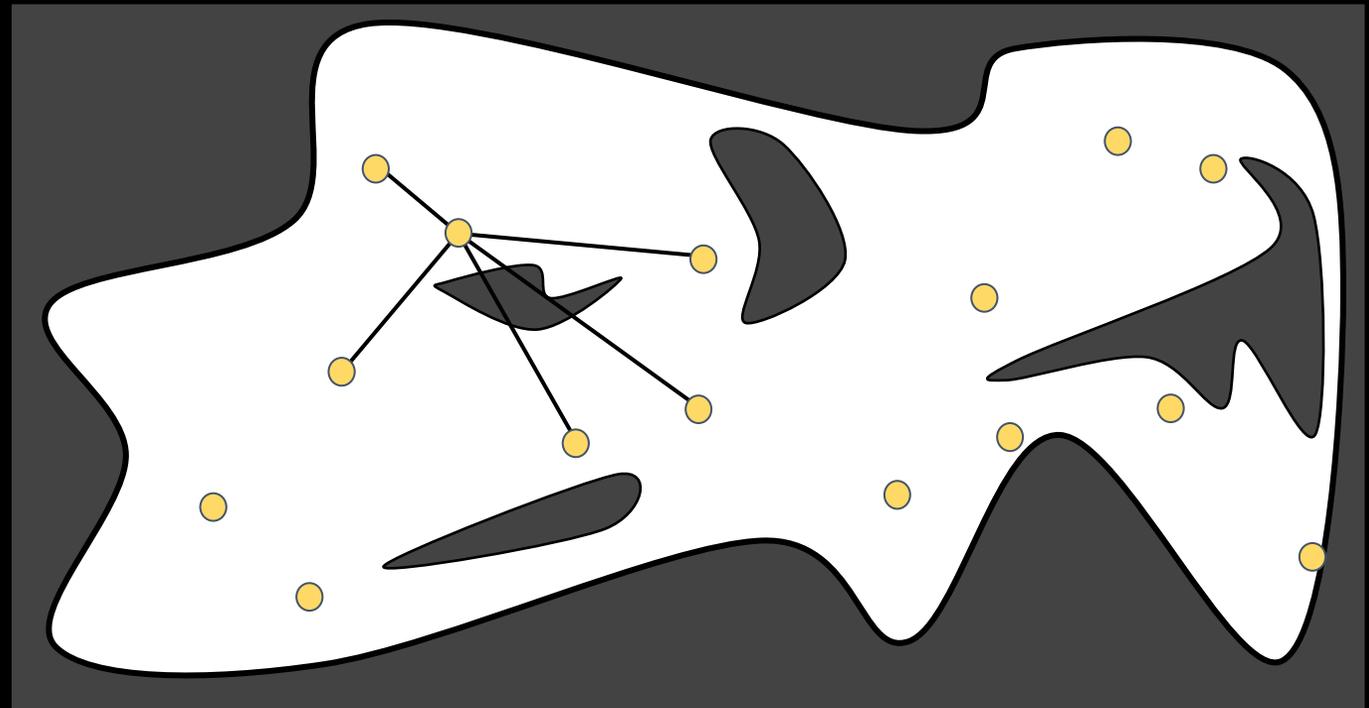
Probabilistic Roadmaps

- Configurations are sampled by picking coordinates at random
- Sampled configurations are tested for collision



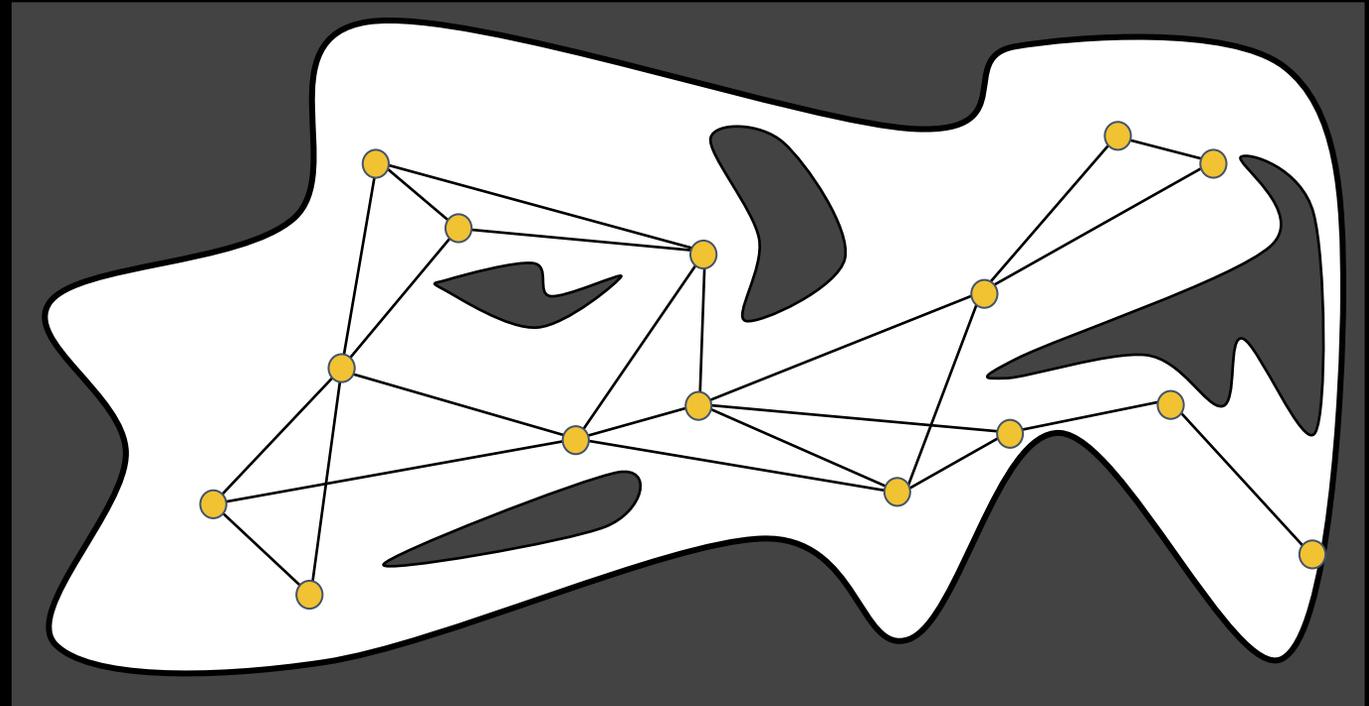
Probabilistic Roadmaps

- Configurations are sampled by picking coordinates at random
- Sampled configurations are tested for collision
- Each configuration is linked by straight paths to its nearest neighbors



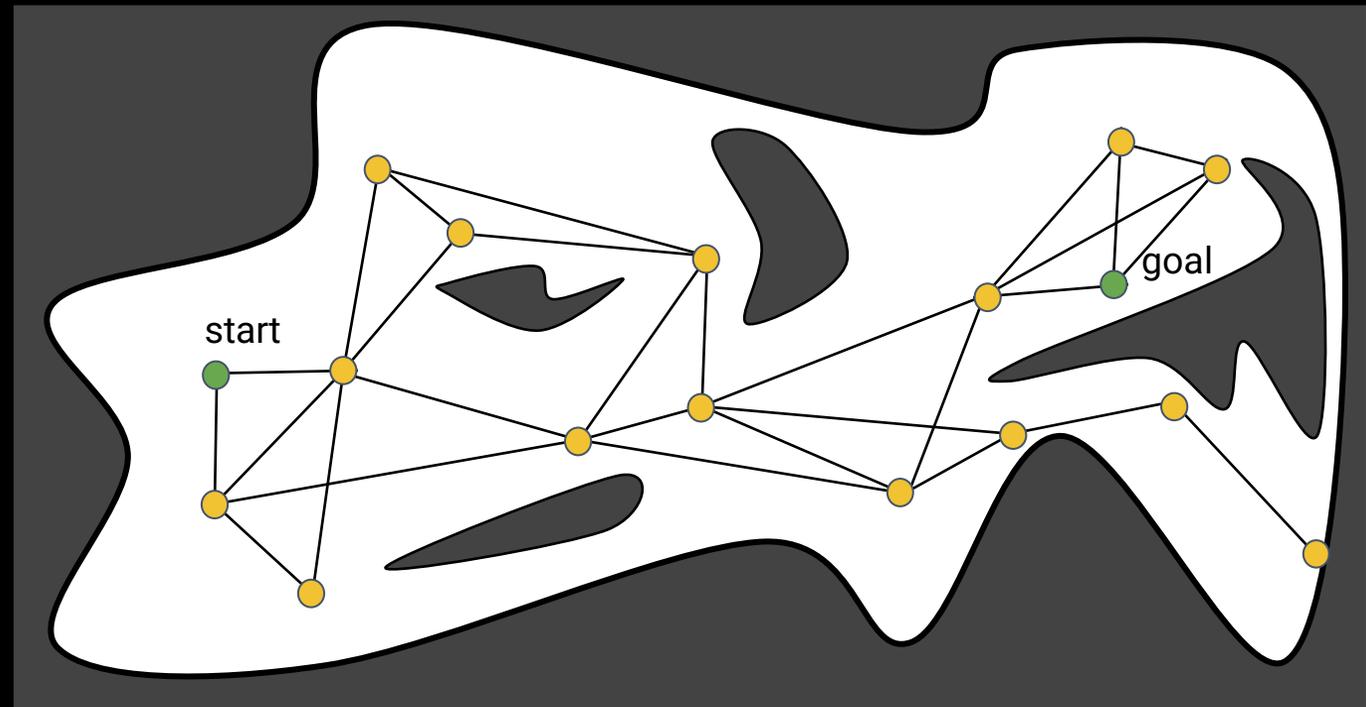
Probabilistic Roadmaps

- Configurations are sampled by picking coordinates at random
- Sampled configurations are tested for collision
- Each configuration is linked by straight paths to its nearest neighbors
- The collision-free links are retained as local paths to form the PRM



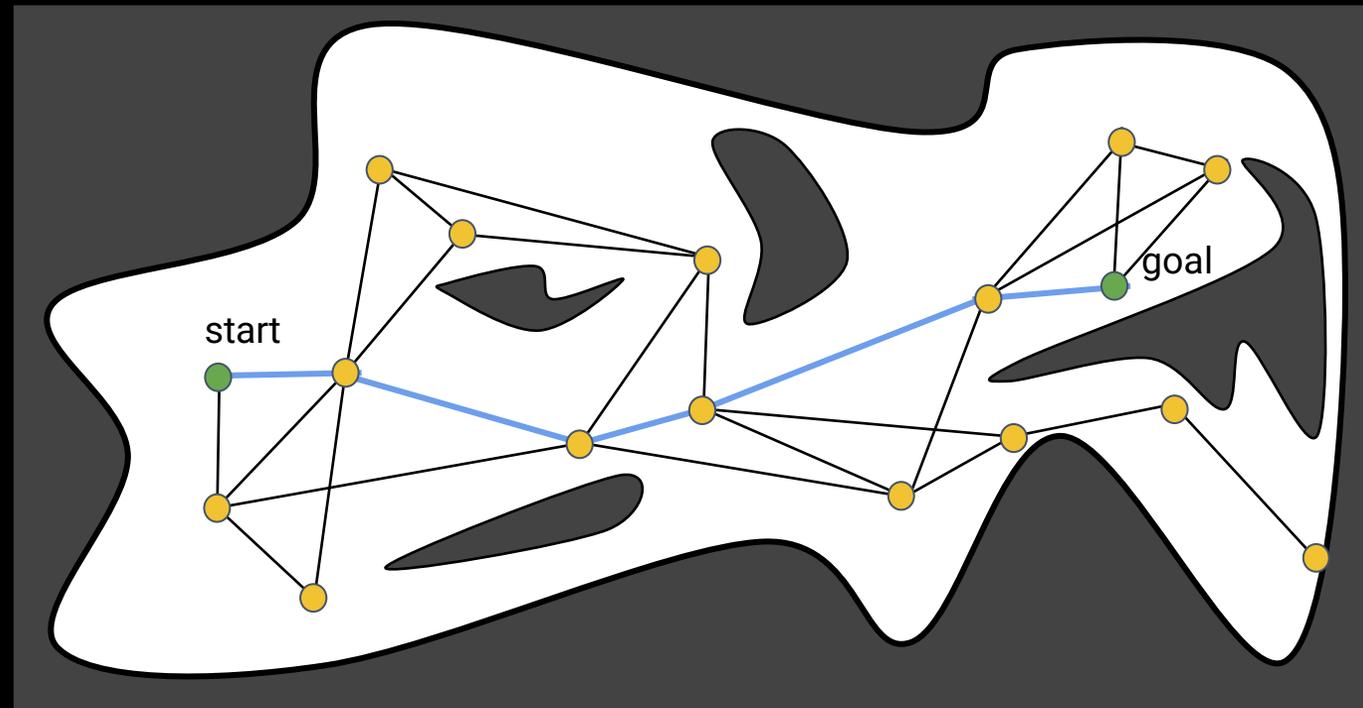
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- The start and goal configurations are included as milestones



Probabilistic Roadmaps

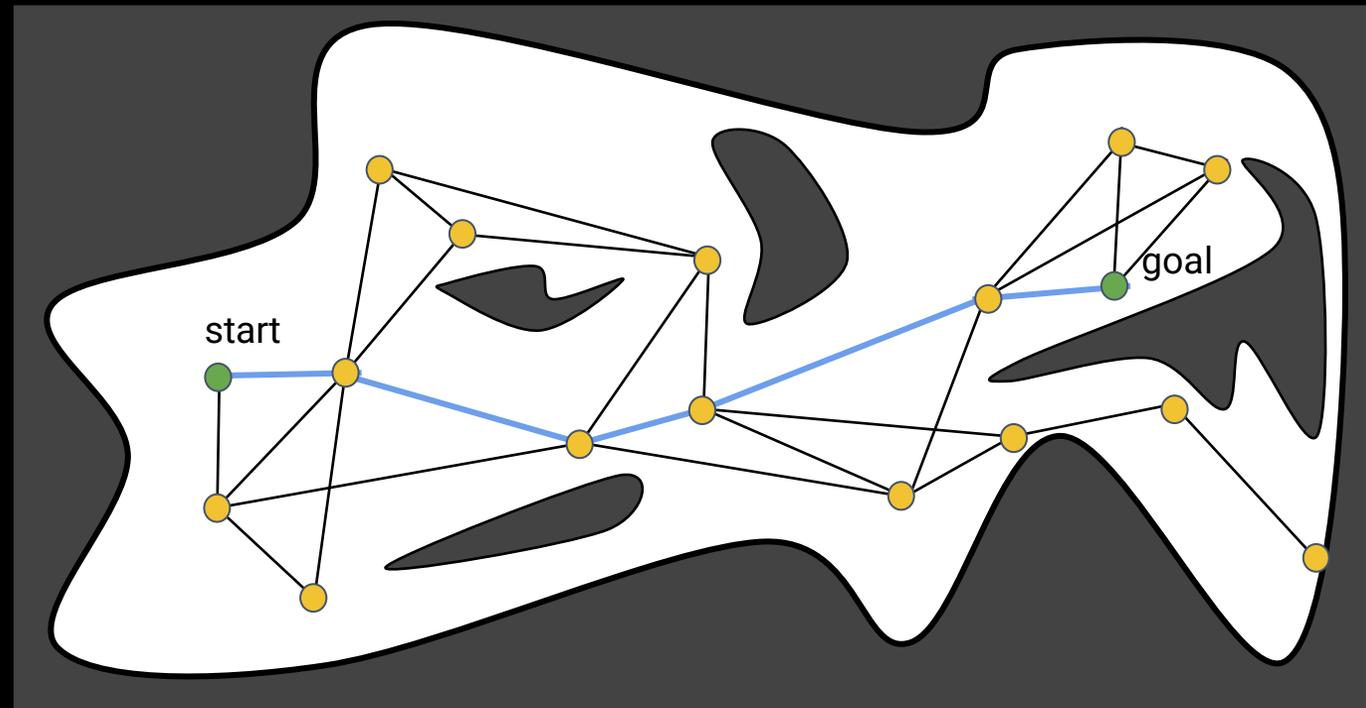
- Configurations are sampled by picking coordinates at random
- Sampled configurations are tested for collision
- Each configuration is linked by straight paths to its nearest neighbors
- The collision-free links are retained as local paths to form the PRM
- The start and goal configurations are included as milestones
- The PRM is searched for a path from start to goal



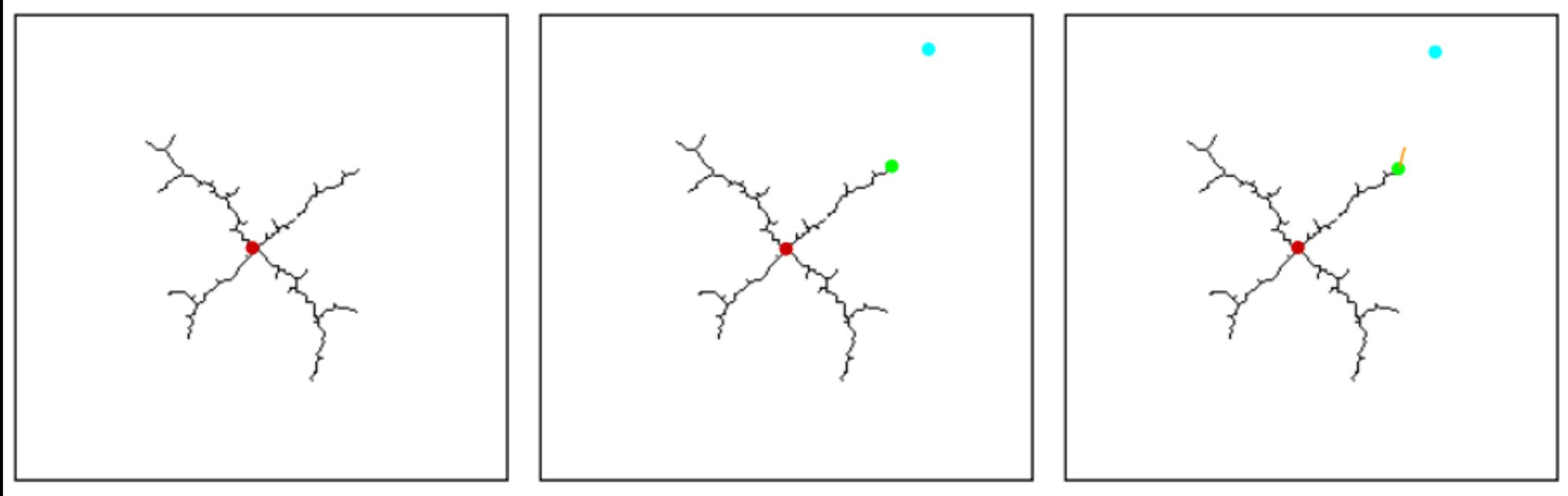
Probabilistic Roadmaps

Considerations

- Single query/multi query
- How are nodes placed?
 - Uniform sampling strategies
 - Non-uniform sampling strategies
- How are local neighbors found?
- How is collision detection performed?
 - Dominates time consumption in PRMs



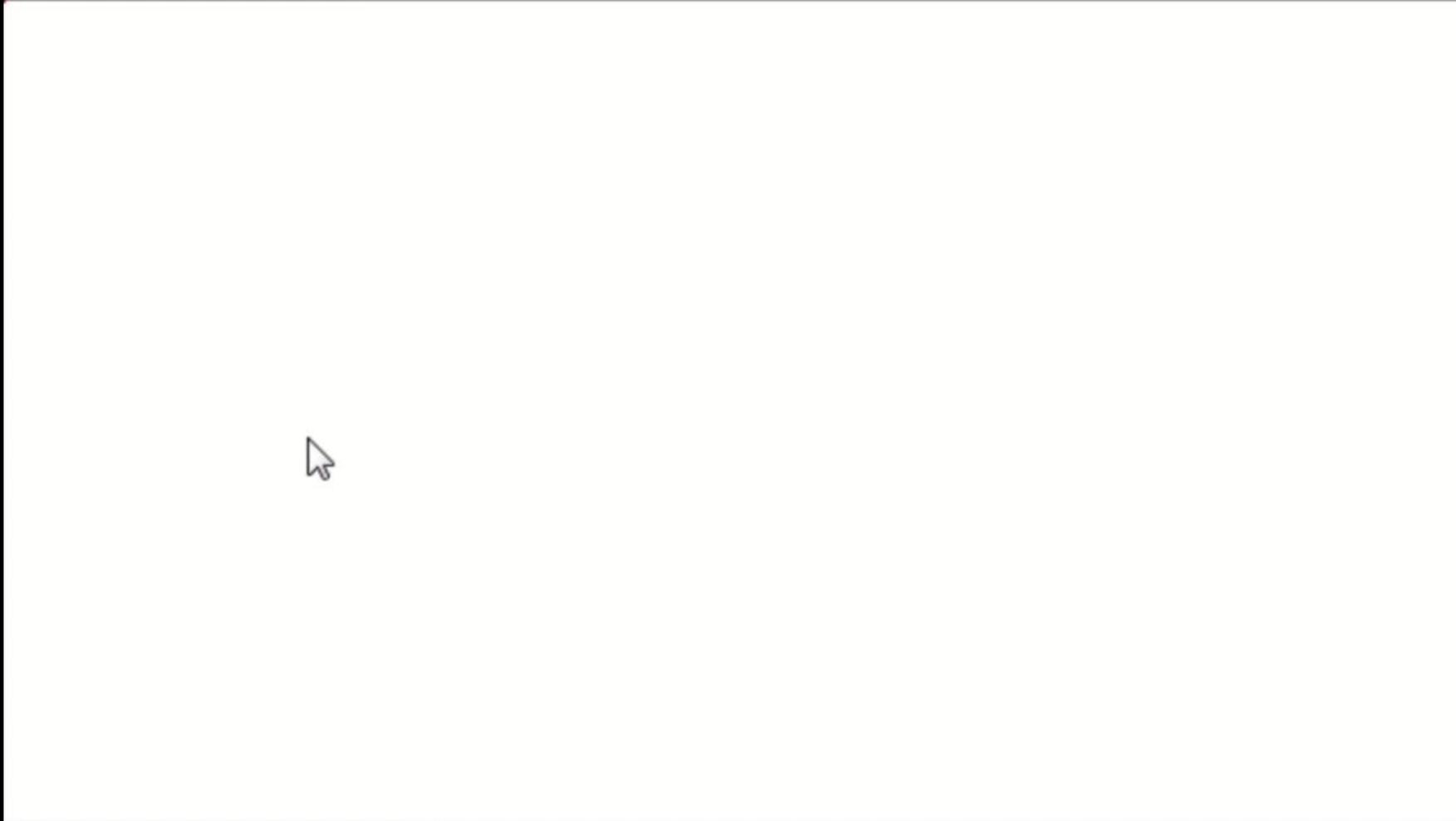
Rapidly Exploring Random Trees (RRT)



1. Maintain a tree rooted at the starting point ●
2. Choose a point at random from free space ●
3. Find the closest configuration already in the tree ●
4. Extend the tree in the direction of the new configuration /

Rapidly Exploring Random Trees (RRT)

- Uniform/biased sampling



Aaron Becker, UH, Wolfram Player example

Rapidly Exploring Random Trees (RRT) - Considerations

- **Sensitive to step-size (Δq)**
 - Small: many nodes, closely spaced, slowing down nearest neighbor computation
 - Large: Increased risk of suboptimal plans / not finding a solution
- **How are samples chosen?**
 - Uniform sampling may need too many samples to find the goal
 - Biased sampling towards goal can ease this problem
- **How are closest neighbors found?**
- **How are local paths generated?**
- **Variations**
 - RRT Connect, A*-RRT, Informed RRT*, Real-Time RRT*, Theta*-RRT, etc.

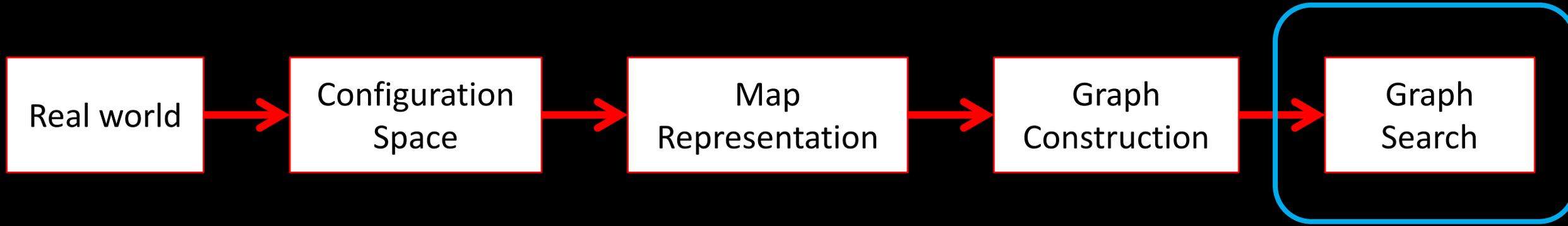
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Fast Robots

Graph Search

Modelling path planning as a graph search problem



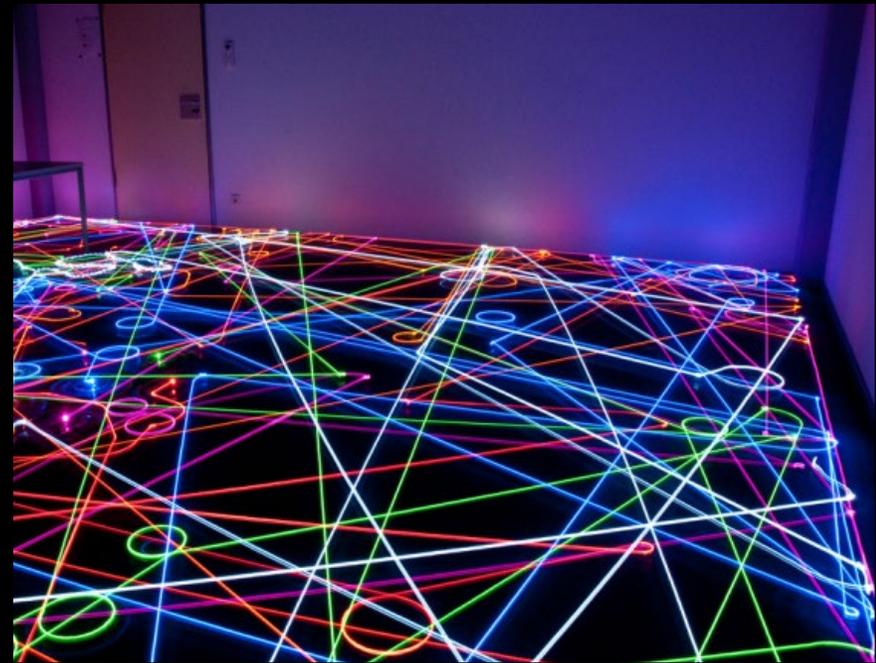
https://pythonrobotics.readthedocs.io/en/latest/modules/path_planning.html#basic-rrt

- Breadth first
- Depth first
- Dijkstra
- A*



Graph Search

- **What is the simplest thing to do?**
 - Random or brute force search
- **Other methods?**
 - Uninformed search
 - Depth First Search (DFS)
 - Breadth First Search (BFS)
 - Dijkstra's Search (LCFS)
 - Informed Search
 - Greedy
 - A*
 - (and many more)



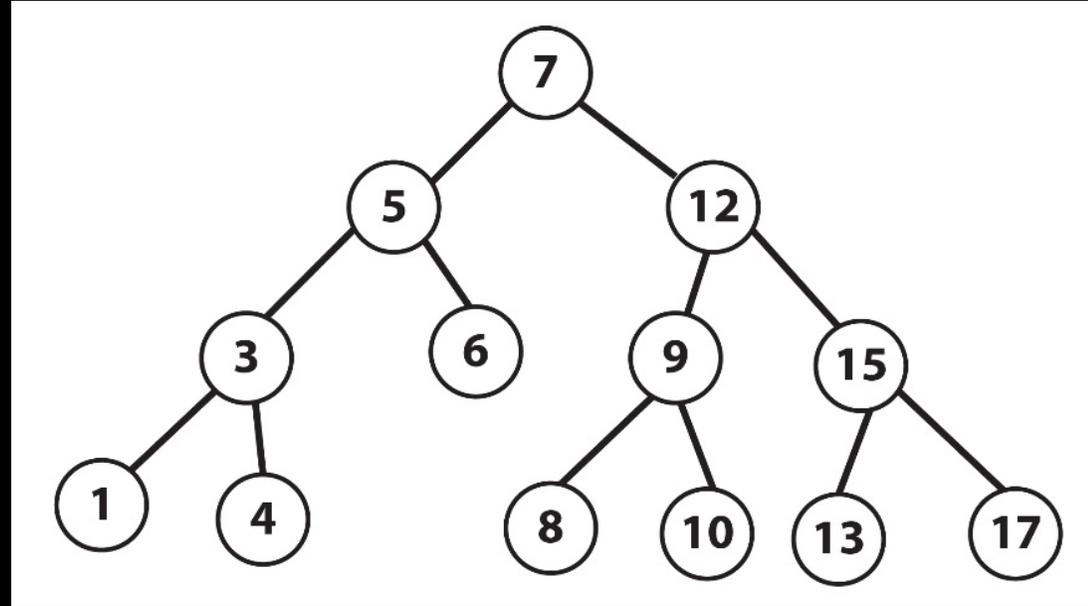
Comparing Search Algorithms

Vocabulary

- Node, edge, parents/children, branching factor, depth

Definitions

- Complete
 - Guaranteed to find a solution in finite time
- Time complexity
 - Worst-case run time
- Space complexity
 - Worst-case memory
- Optimality
 - A search is optimal if it is complete, and only returns cost-minimizing solutions

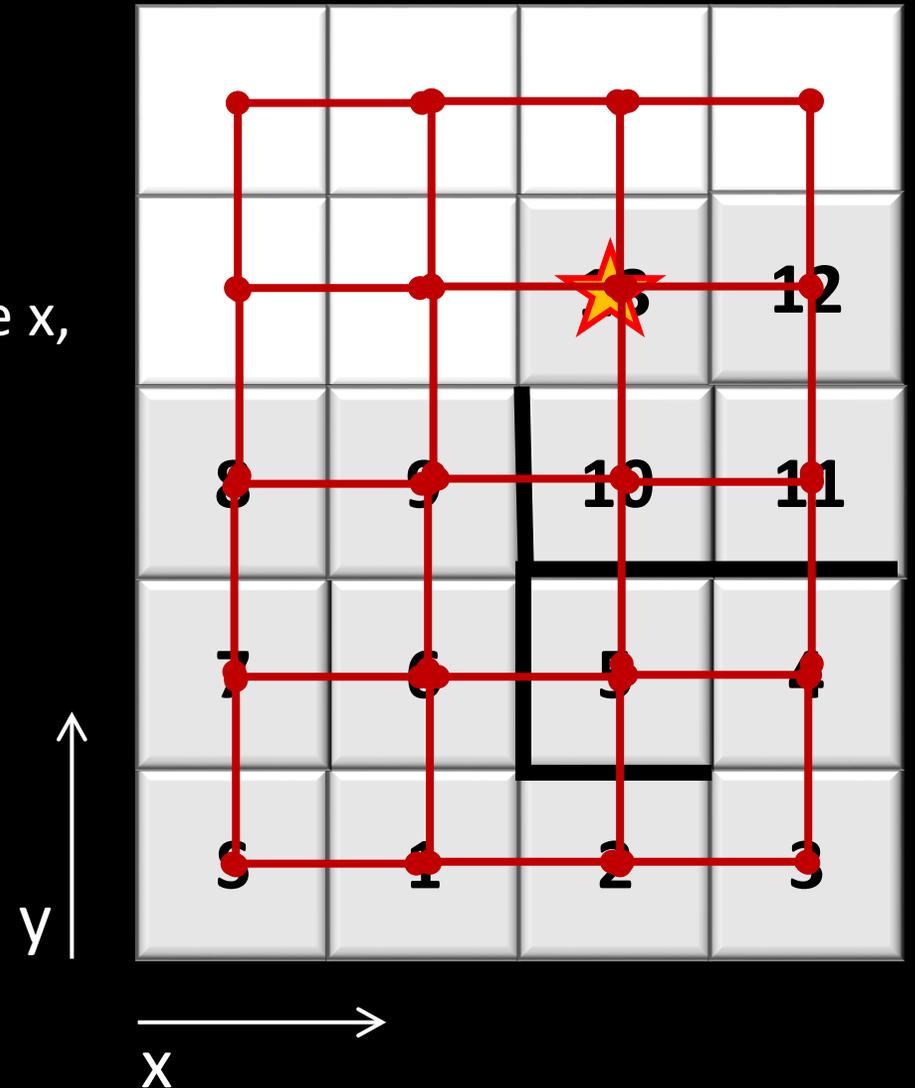


Algorithms and Search

- **What is the simplest thing to do?**
 - Random or brute force search
- **How many grid traversals will brute force take?**
 - First establish a search order
 - Advance x first, then increment y and decrease x, etc.

Search order: N, E, S, W

Find a goal

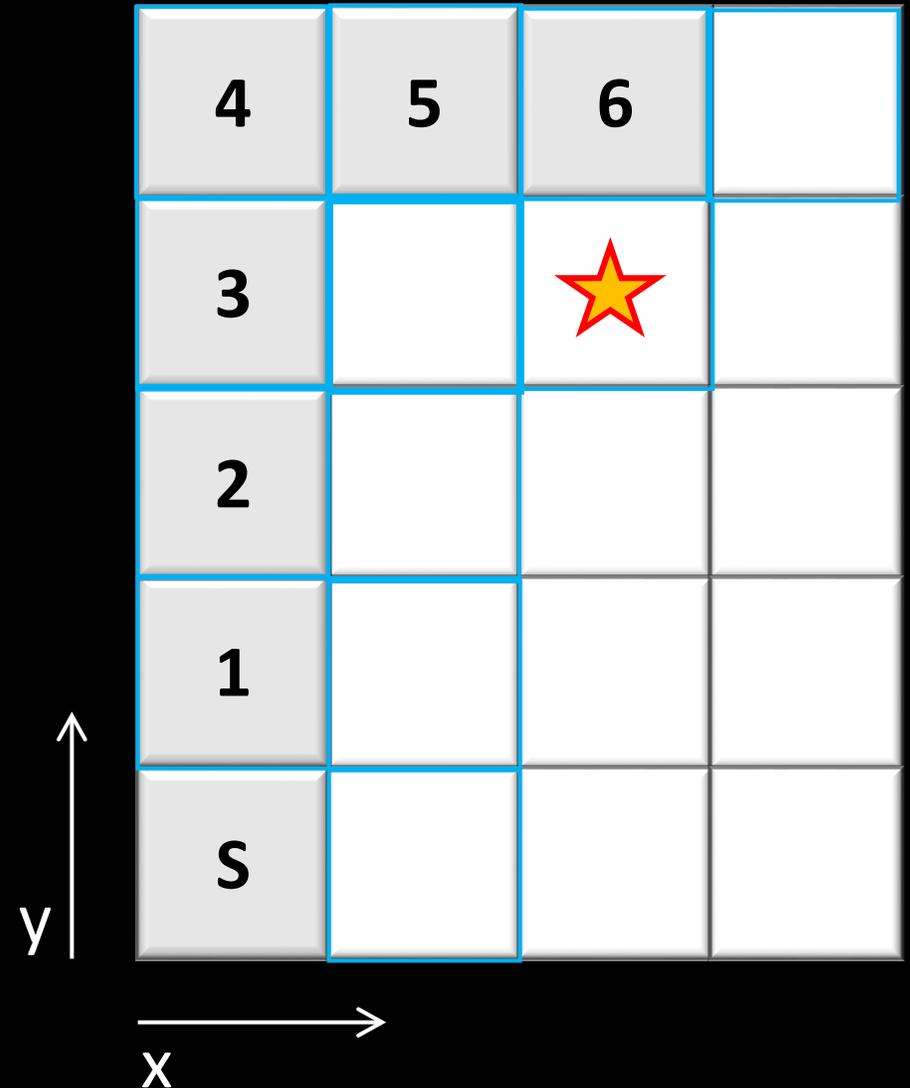


Algorithms and Search

- **What is the simplest thing to do?**
 - Random or brute force search
- **Other methods?**
 - Depth First Search (DFS)

Search order: N, E, S, W

Find a goal

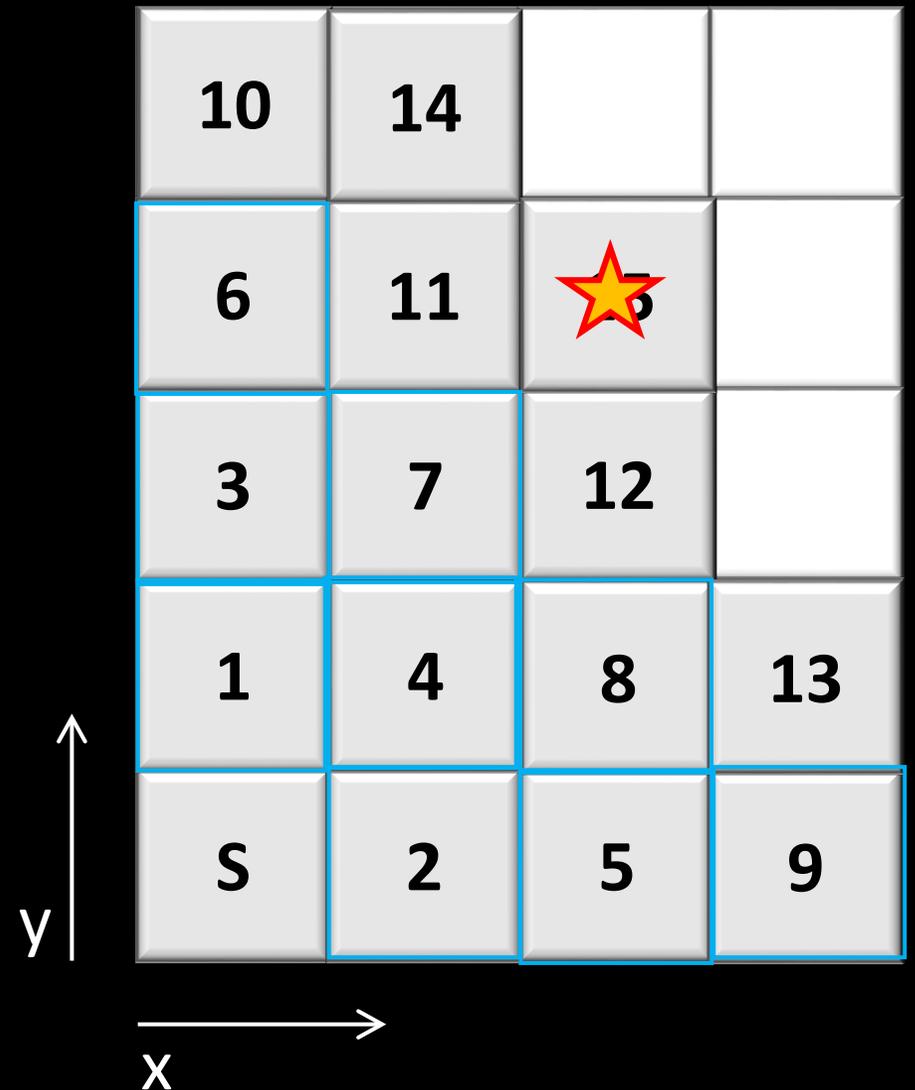


Algorithms and Search

- **What is the simplest thing to do?**
 - Random or brute force search
- **Other methods?**
 - Depth First Search (DFS)
 - Breadth First Search (BFS)

Search order: N, E, S, W

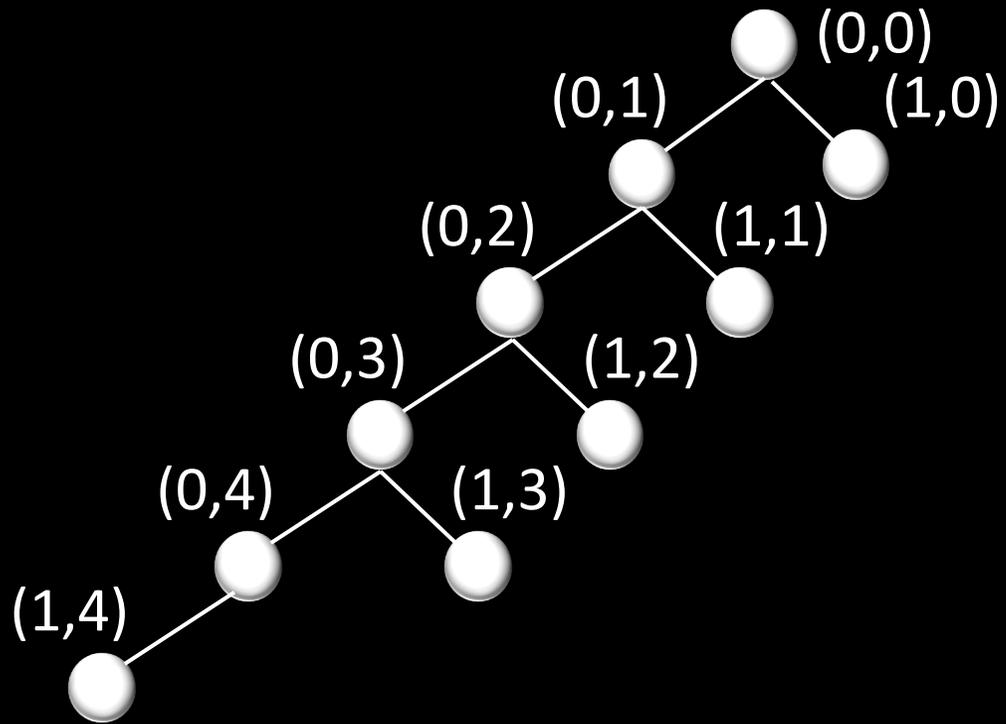
Find a goal



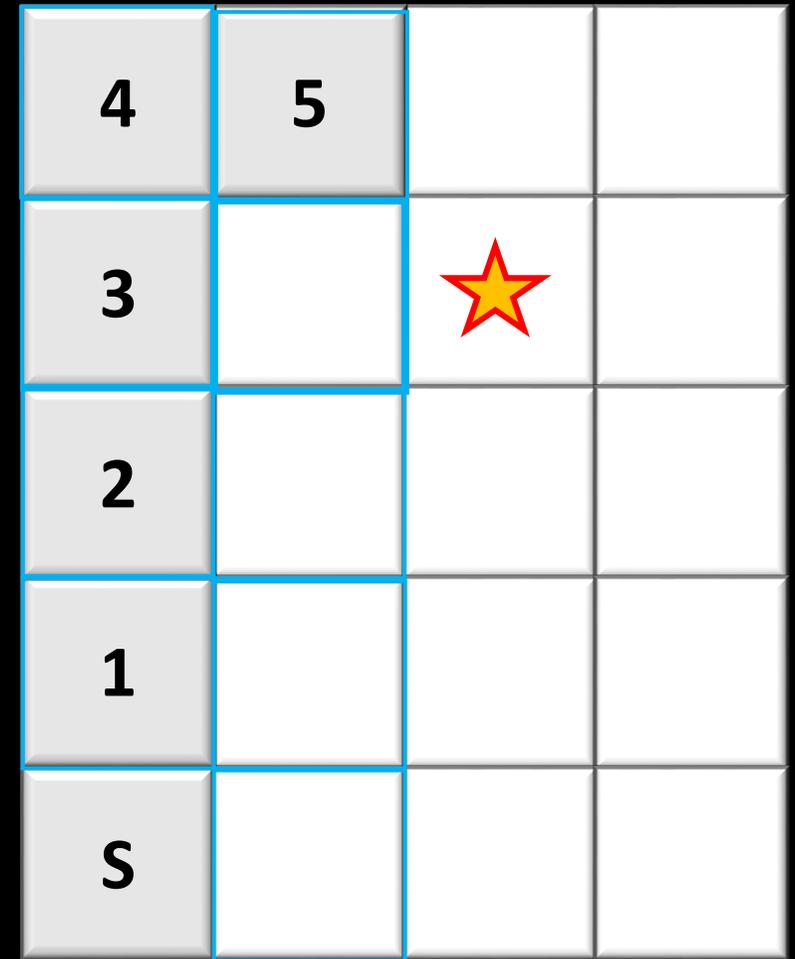
Depth First Search (DFS)

Search order: N, E, S, W

Find a goal



and so on...



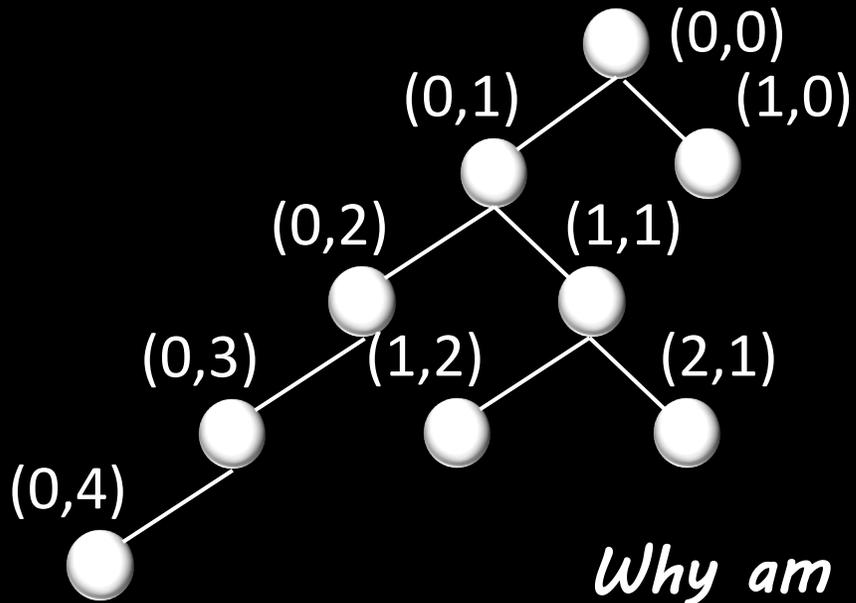
y ↑

→ x

Depth First Search (DFS)

Search order: N, E, S, W

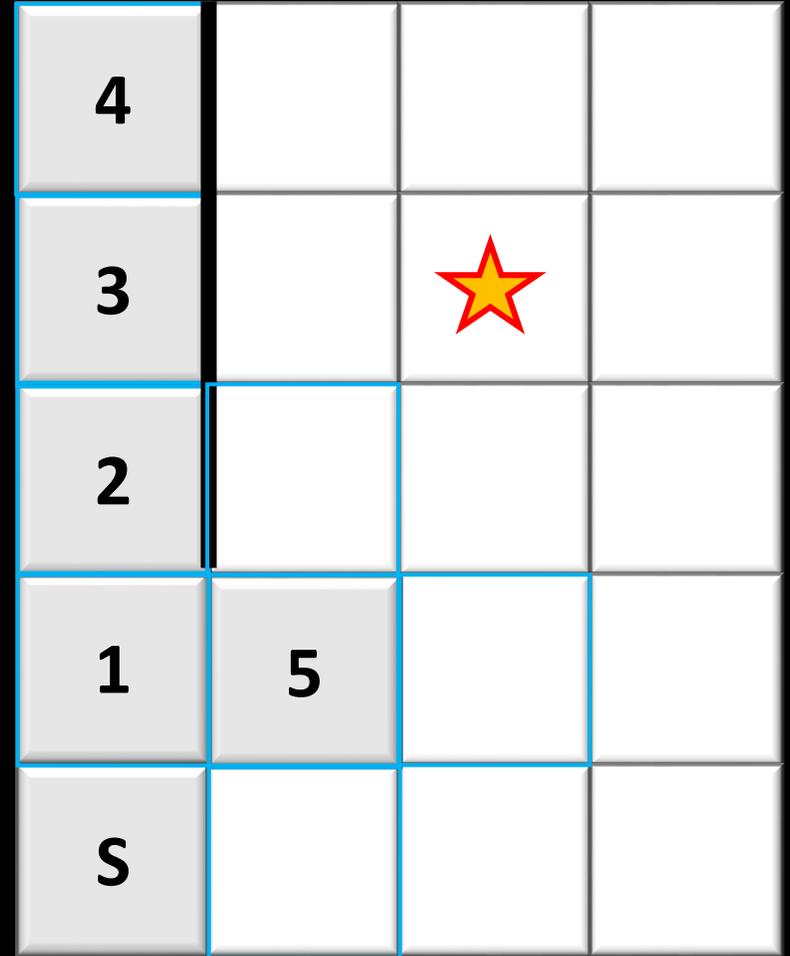
Find a goal



Why am I not also adding (1,0)?

- Keep track of what is already on the frontier

and so on...



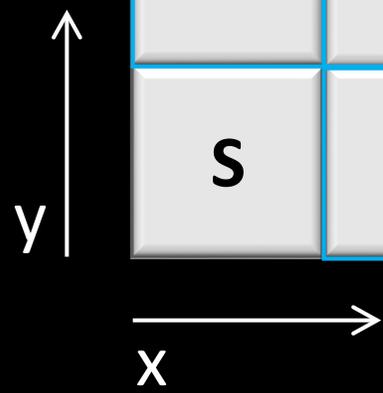
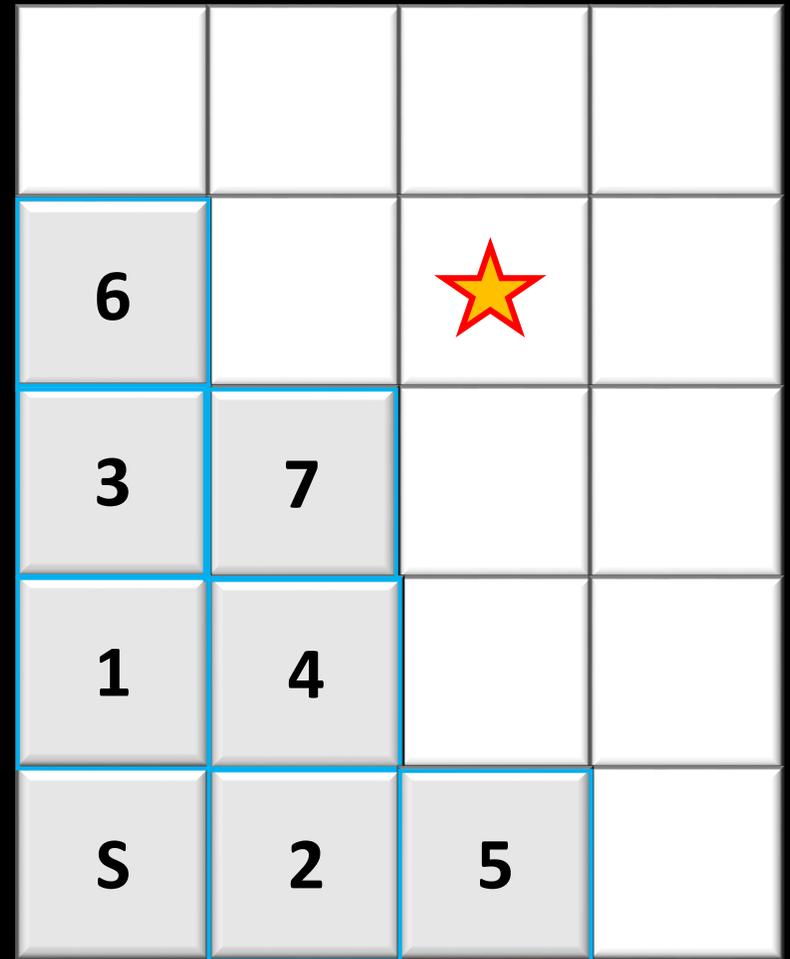
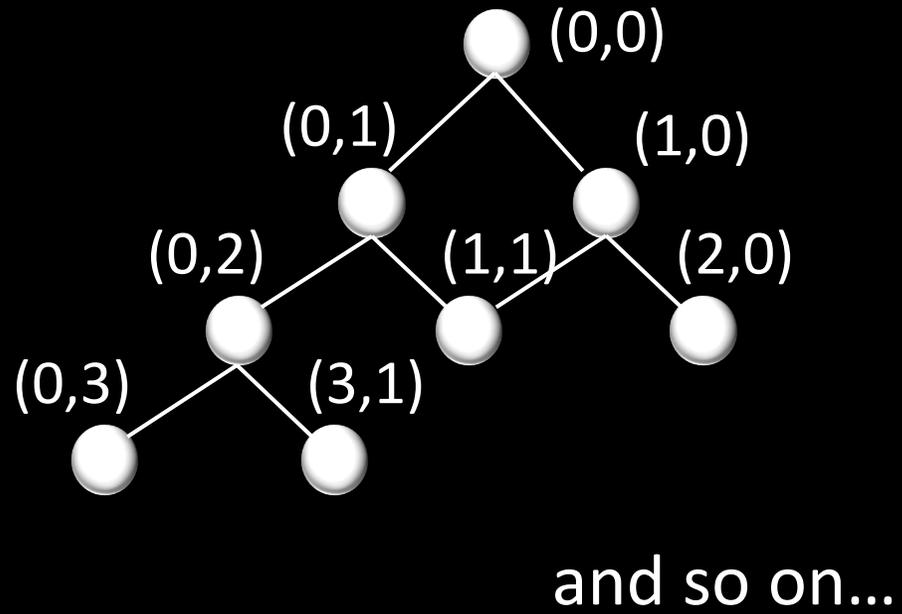
y ↑

→ x

Breadth First Search (BFS)

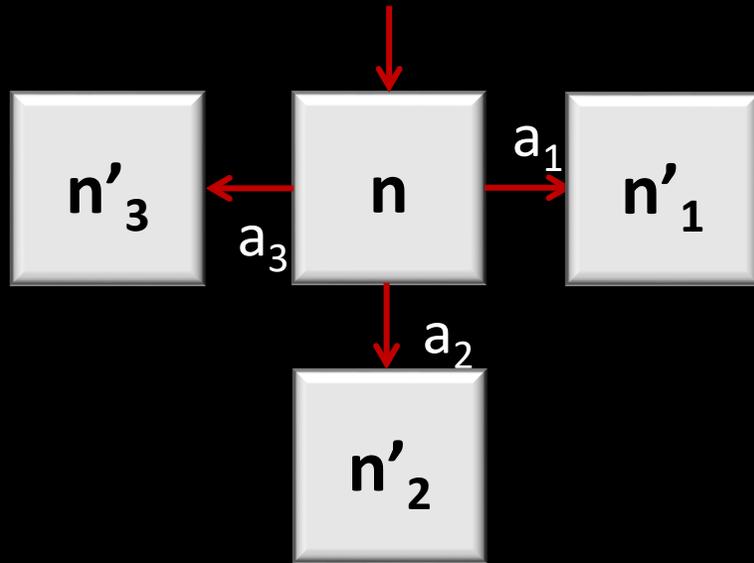
Search order: N, E, S, W

Find a goal

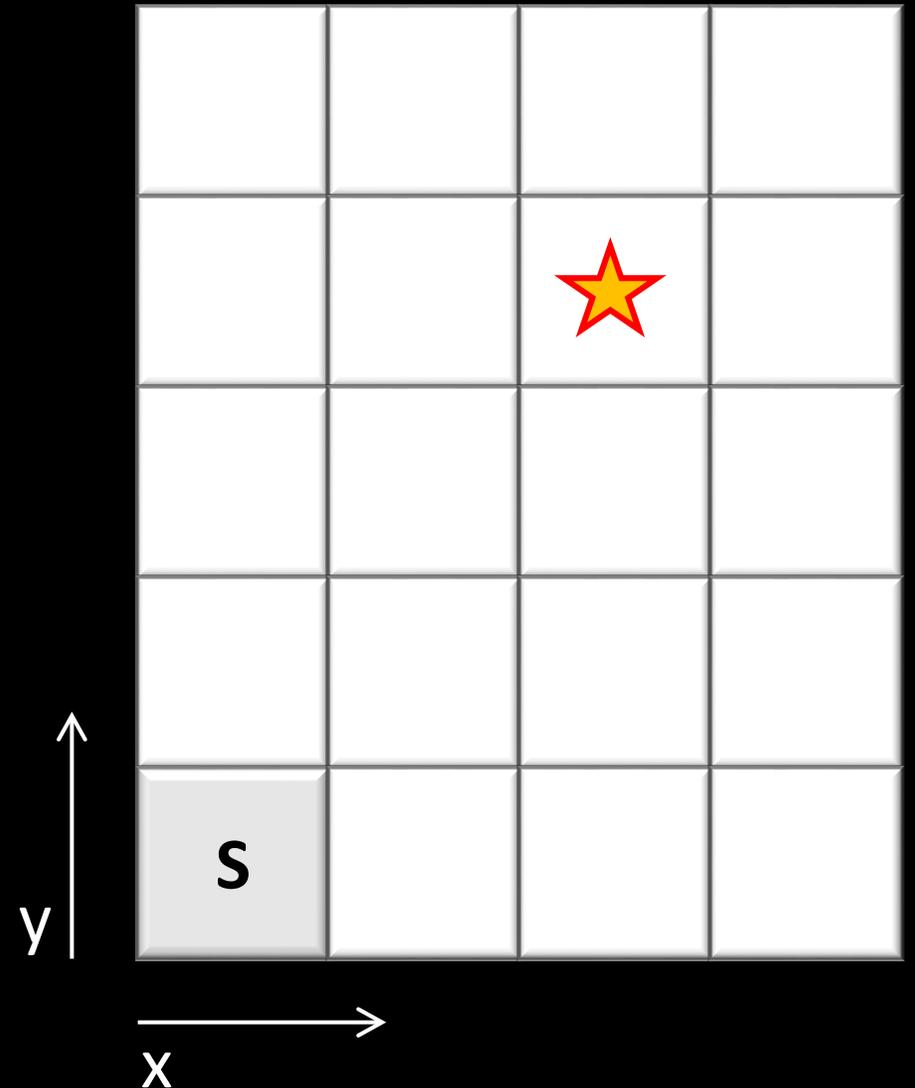


Search Algorithms, General

- Common graph structure
 - For every node, n
 - you have a set of actions, a
 - that moves you to a new node, n'



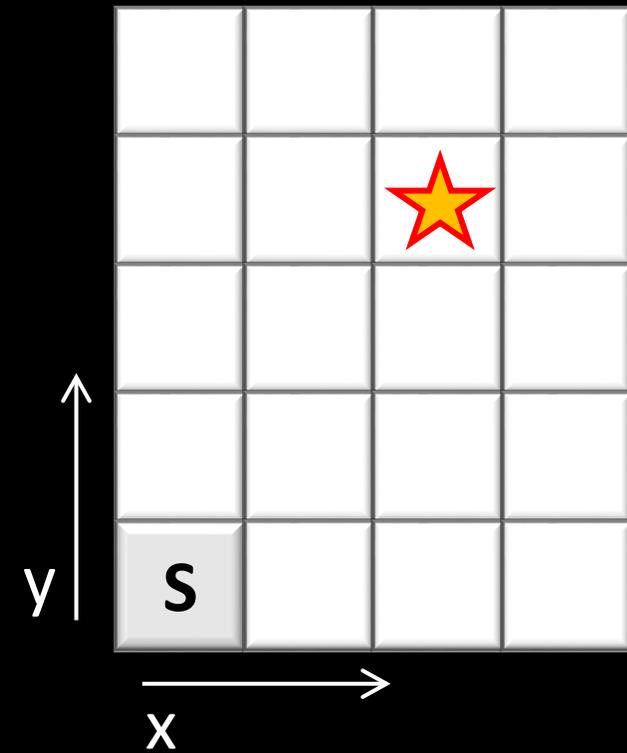
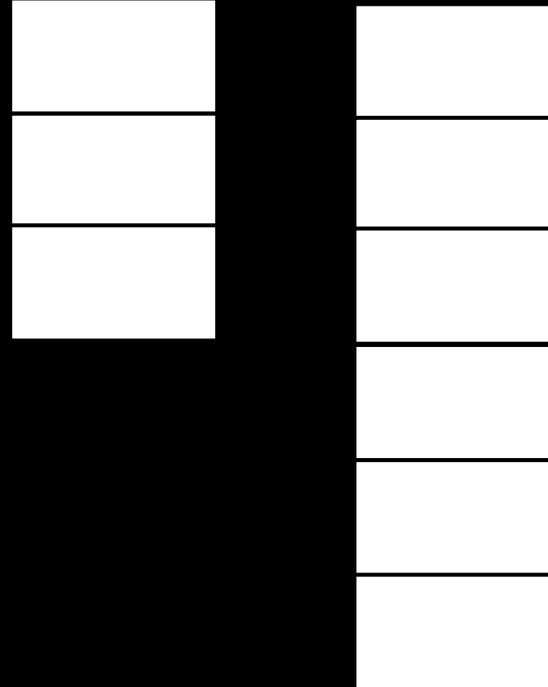
Find a goal



Search Algorithms, General

```
n = state(init)
frontier.append(n)
while(frontier not empty)
  n = pull state from frontier
  append n to visited
  if n = goal, return solution
  for all actions in n
    n' = a(n)
    if n' not visited
      append n' to frontier
```

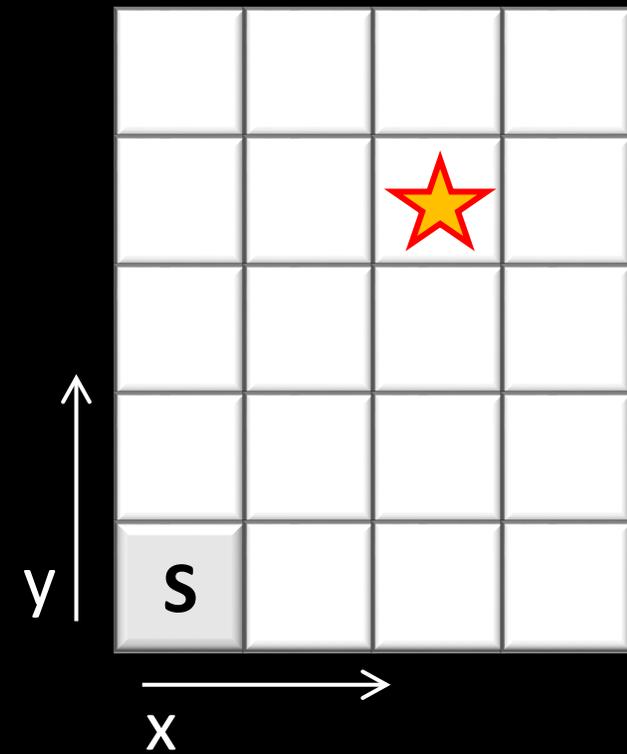
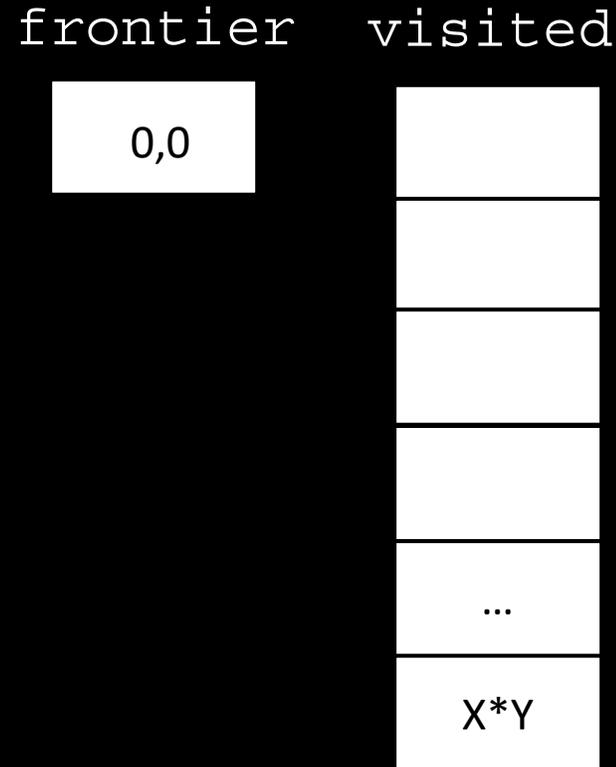
frontier visited



*How much space
to allocate to
your buffers?*

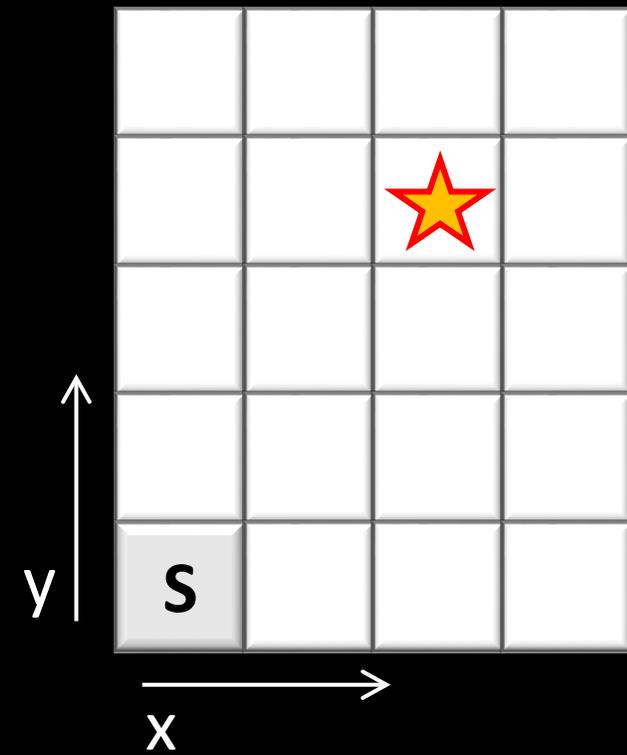
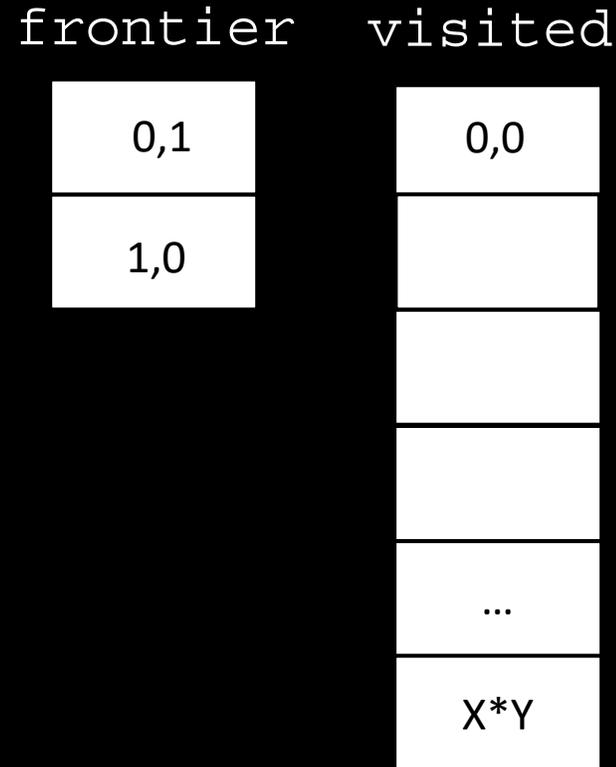
Depth First Search (DFS)

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Depth First Search (DFS)

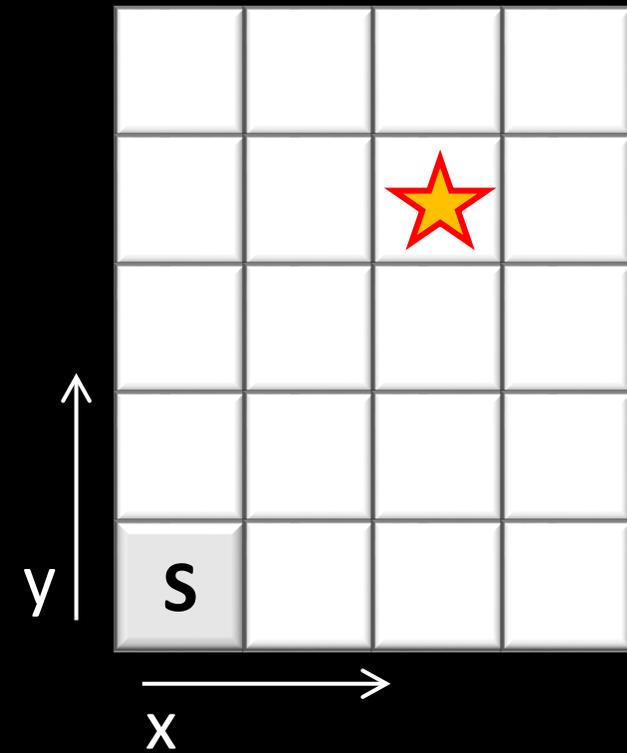
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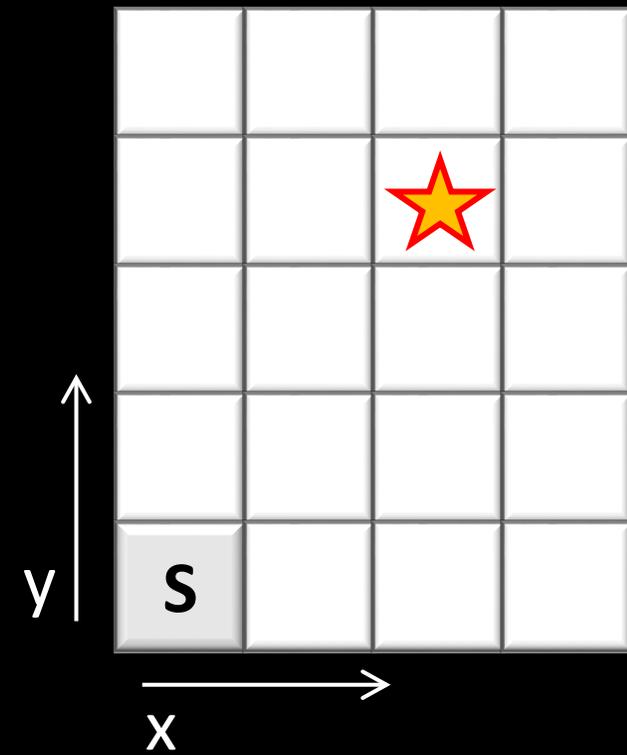
frontier	visited
0,1	0,0
1,0	0,1
1,0	
	...
	X*Y



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      append n' to frontier
```

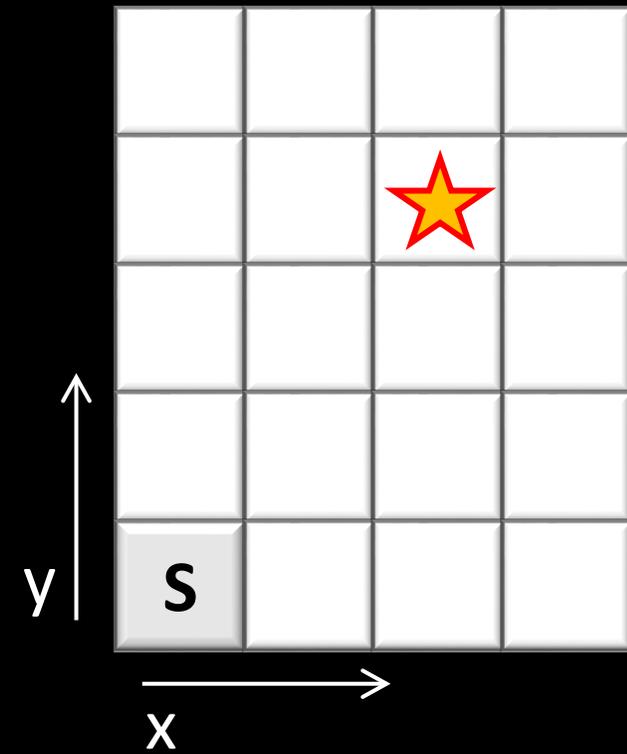
frontier	visited
0,2	0,0
1,1	0,1
1,0	0,2
	...
	X*Y



Depth First Search (DFS)

```
n = state(init)
frontier.append(n)
while(frontier not empty)
  n = pull state from frontier
  append n to visited
  if n = goal, return solution
  for all actions in n
    n' = a(n)
    if n' not visited
      append n' to frontier
```

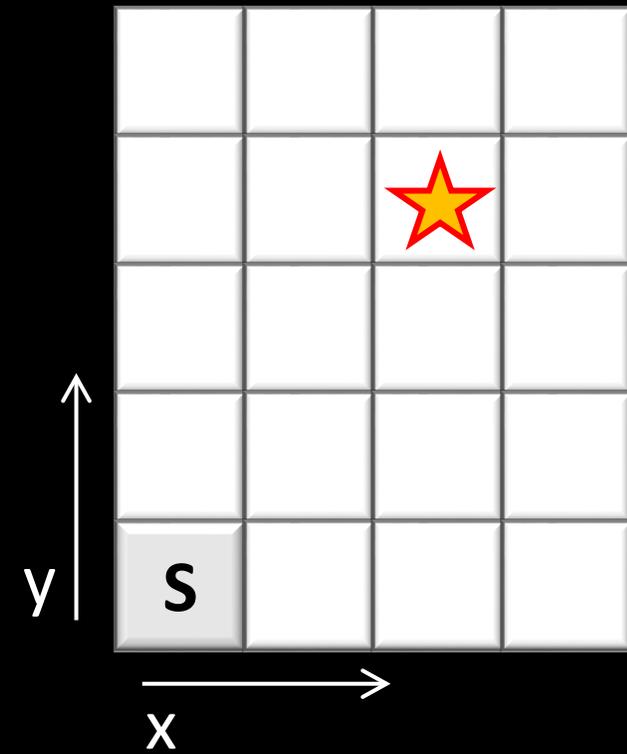
frontier	visited
0,3	0,0
1,2	0,1
1,0	0,2
1,0	0,3



Depth First Search (DFS)

```
n = state(init)
frontier.append(n)
while(frontier not empty)
  n = pull state from frontier
  append n to visited
  if n = goal, return solution
  for all actions in n
    n' = a(n)
    if n' not visited
      append n' to frontier
```

frontier	visited
0,4	0,0
1,2	0,1
1,1	0,2
1,0	0,3
1,0	
etc...	

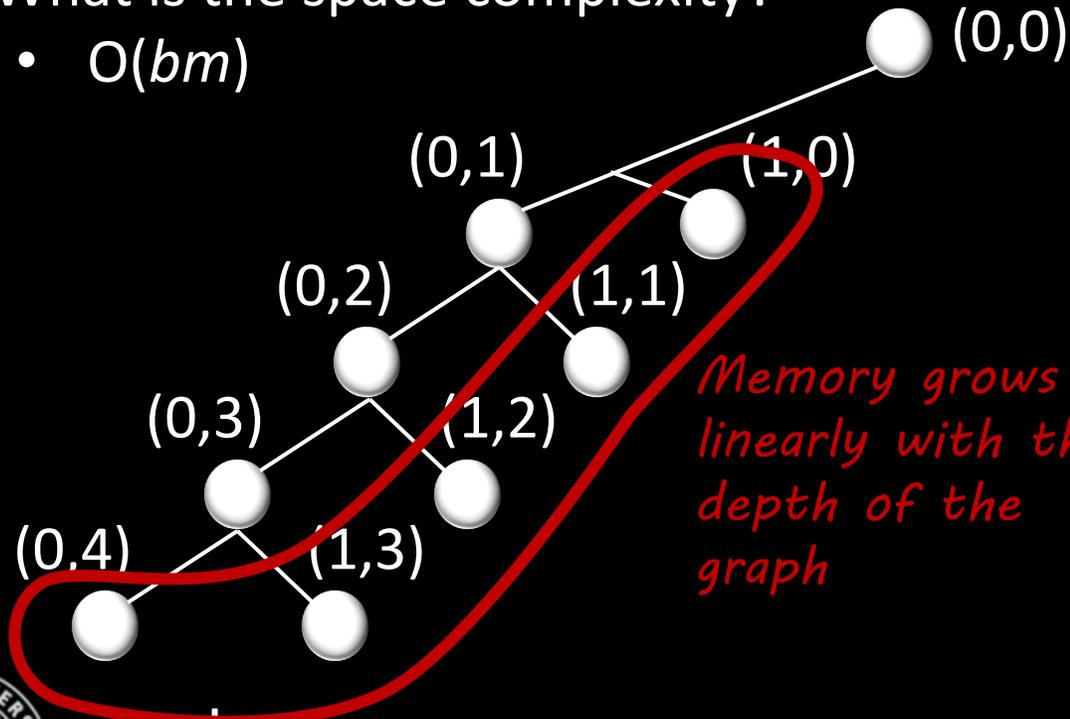


Frontier Buffer?

- *Last-In First-Out (LIFO) Buffer*

Depth First Search (DFS)

- Is it complete?
 - Yes, but only on finite graphs
- What is the time complexity?
 - $O(b^m)$
- What is the space complexity?
 - $O(bm)$



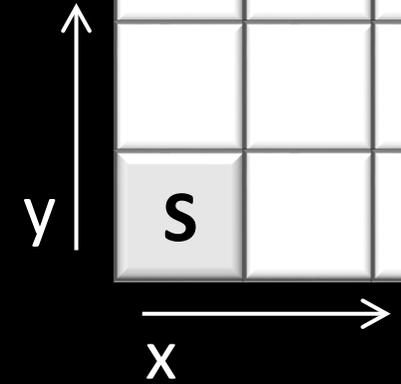
Memory grows linearly with the depth of the graph

and so on...

frontier visited

0,4
1,3
1,2
1,1
1,0

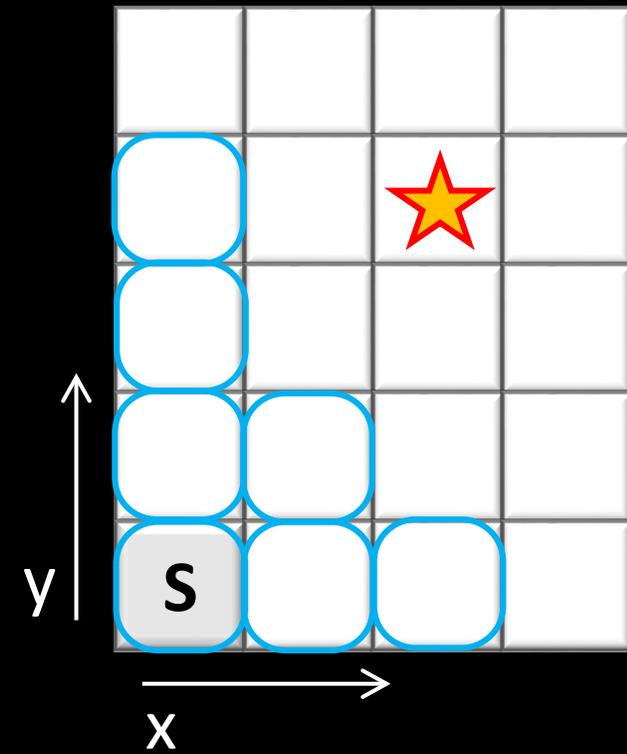
0,0
0,1
0,2
0,3
...
X*Y



Breadth First Search (BFS)

```
n = state(init)
frontier.append(n)
while(frontier not empty)
  n = pull state from frontier
  if n is goal, return solution
  for all actions in n
    n' = a(n)
    if n' not visited
      append n' to frontier
```

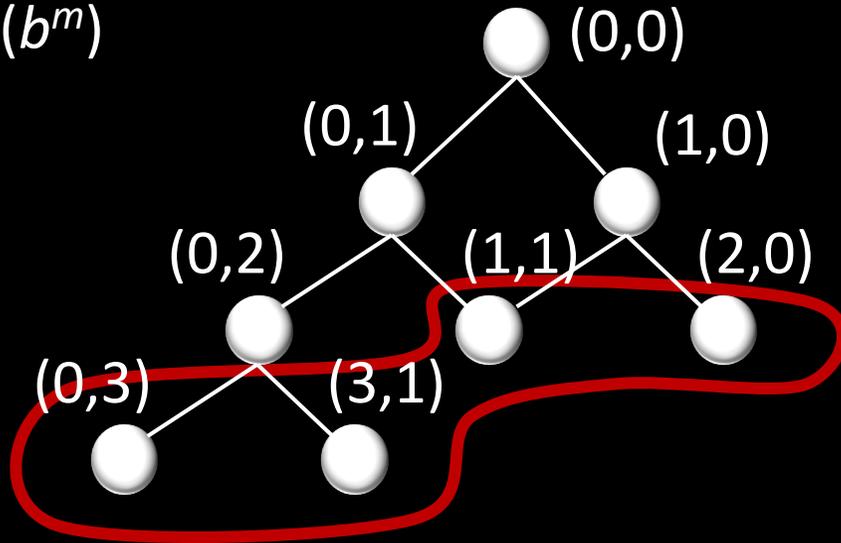
frontier	visited
0,0	0,0
0,1	0,1
1,0	1,0
0,2	0,2
1,1	
2,0	...
0,3	



Type of Buffer?
First-In First-Out (FIFO) Buffer

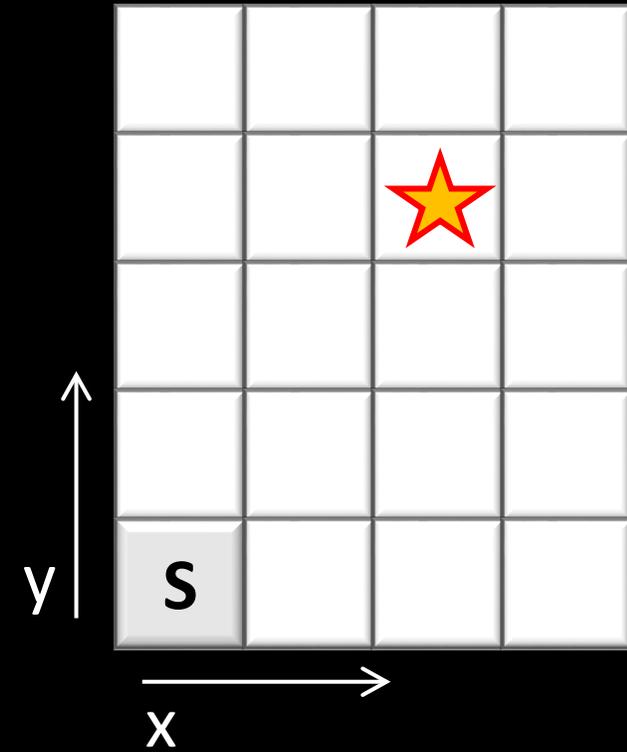
Breadth First Search (BFS)

- Is it complete?
 - Yes, as long as b is finite
- Is it optimal?
 - Yes
- What is the time complexity?
 - $O(b^m)$
- What is the space complexity?
 - $O(b^m)$



frontier visited

	0,0
	0,1
	1,0
	0,2
1,1	
2,0	
0,3	
	...



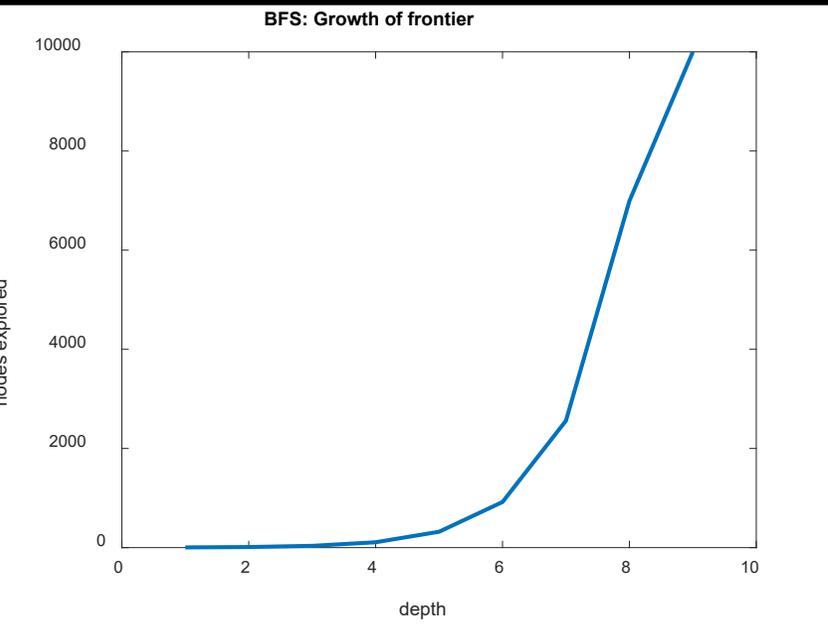
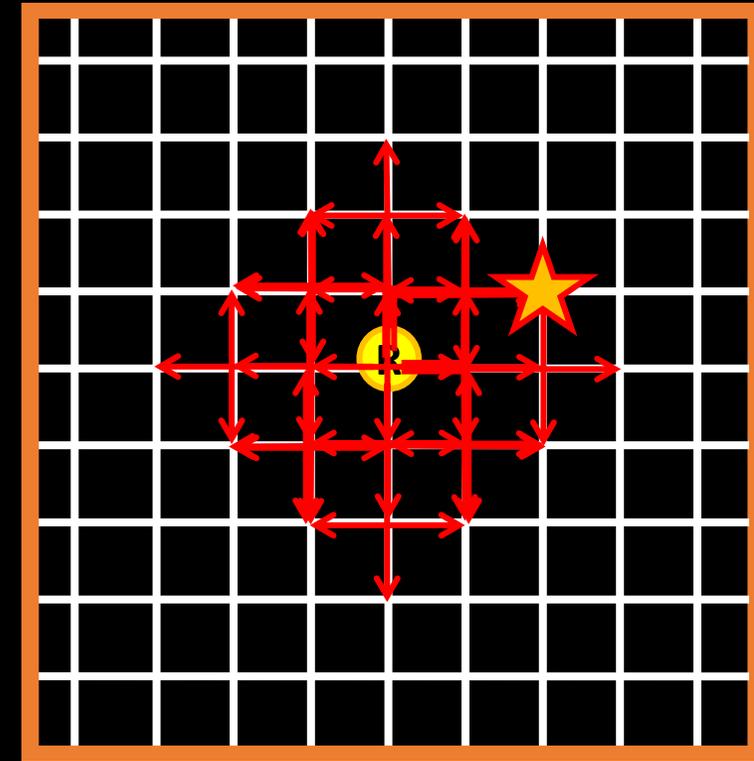
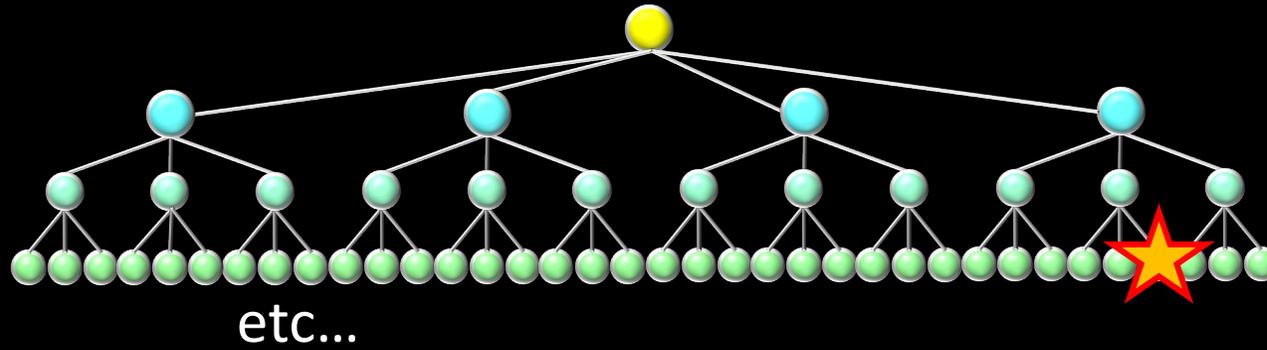
*Type of Buffer?
First-In First-Out (FIFO) Buffer*

*Memory grows exponentially
with the depth of the graph*

BFS: Memory and Computation

Frontier size:

- 4
- 12
- 36

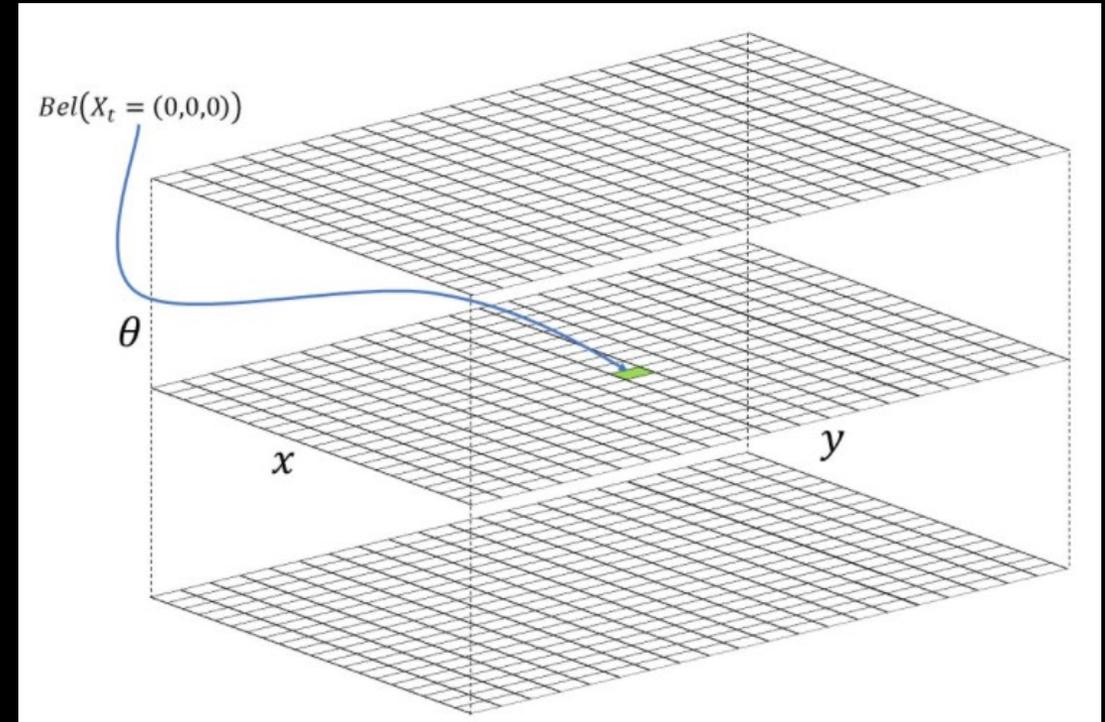


Uninformed Search Algorithms, General

- When is DFS appropriate?
 - If the memory is restricted
 - If solutions tend to occur at the same depth in the tree
- When is DFS inappropriate?
 - If some paths have infinite length / if the graph contains cycles
 - If some solutions are very deep, while others are very shallow
- When is BFS appropriate?
 - If you need to find the shortest path
 - If memory is not a problem
 - If some solutions are shallow
 - If there might be infinite paths
- When is BFS inappropriate?
 - If memory is limited / if the branching factor is very large
 - If solutions tend to be located deep in the tree

ECE4160/5160 – MAE 4910/5910 Fast Robots

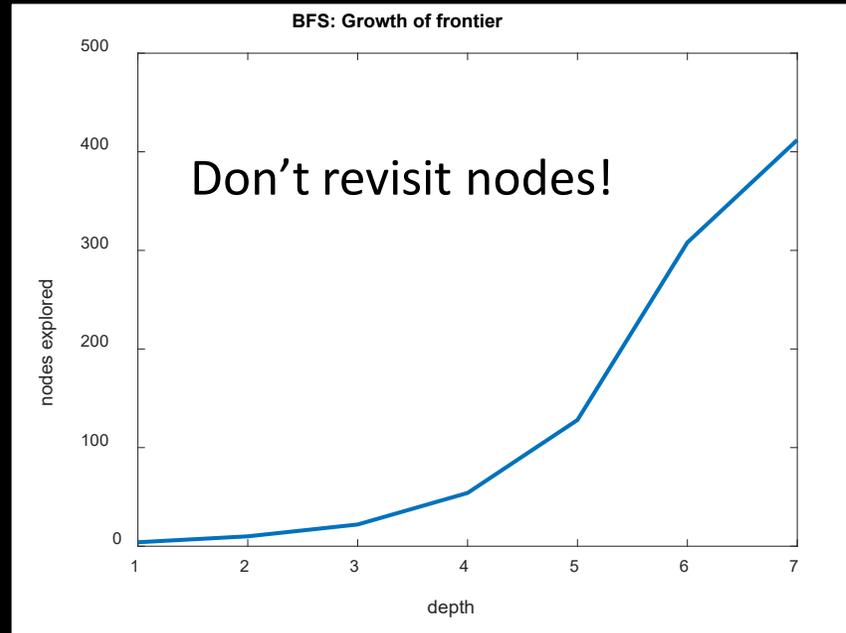
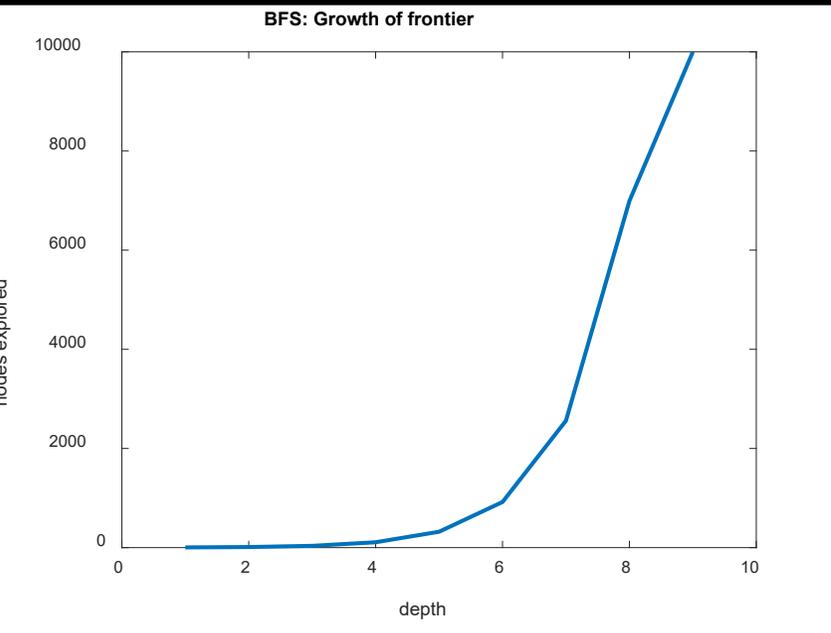
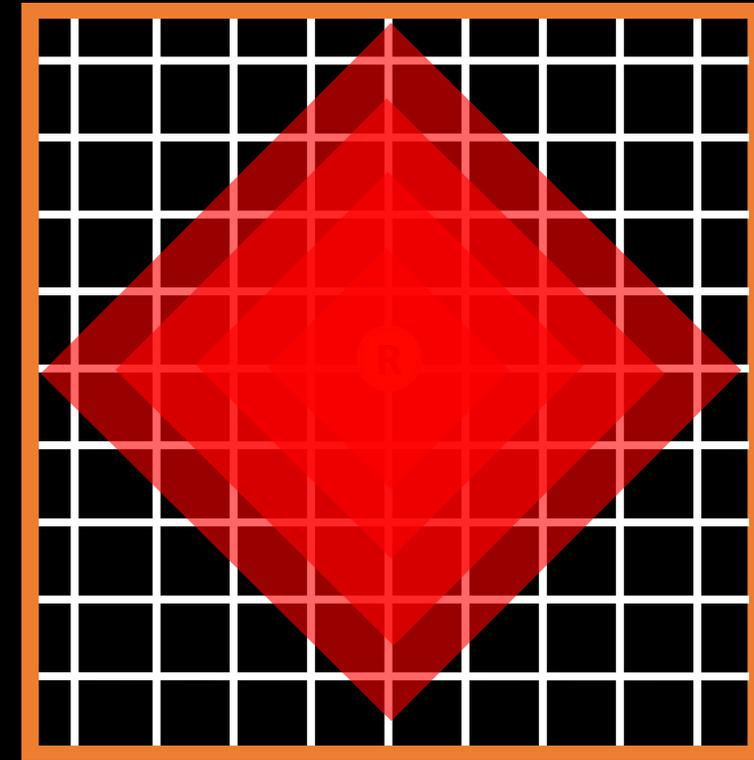
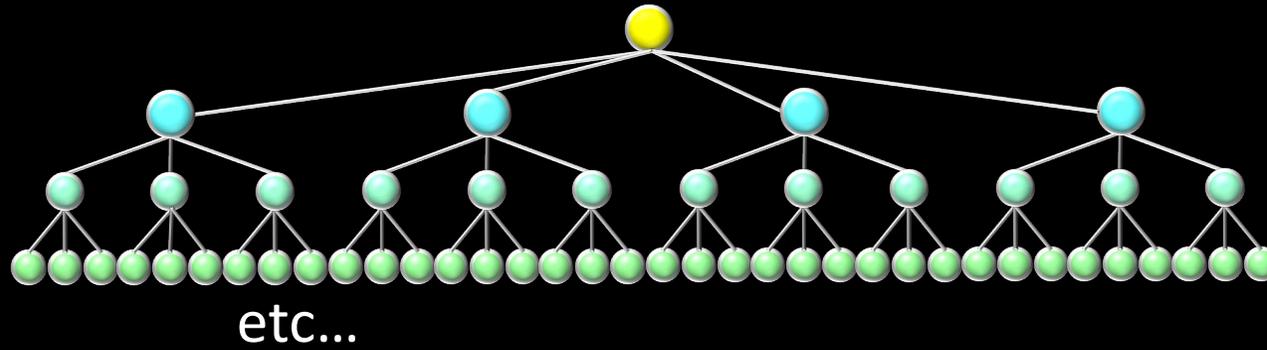
- Is BFS / DFS possible for your task on the Artemis?
 - What is the maximum branching factor?
 - $b = 4$
 - What is the longest path?
 - $m \sim 20 * 20 = 400$
 - Depth First Search
 - Frontier: $O(bm) = 1,600$ nodes
 - Float -> 6.4kB
 - Artemis memory?
 - 1MB flash and 384k RAM
 - Breadth First Search
 - Frontier: $O(b^m) = 4^{20*20} = 6.7e240$ nodes



BFS: Memory and Computation

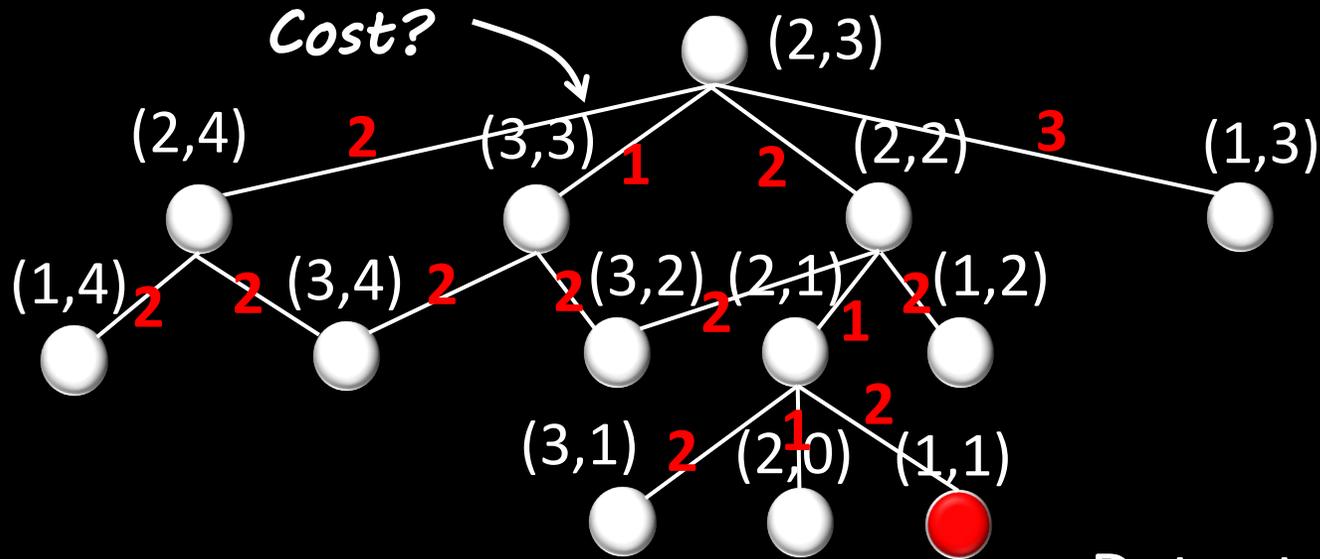
Frontier size:

- 4
- 12
- 36



Lowest-Cost First Search

- Consider parent cost!



What node to expand next?

Data structure

- n.state
- n.parent
- n.cost
- n.action

What cost heuristic could we add?

- Go straight, cost 1
- Turn one quadrant, cost 1

	(1,4)	(2,4)	(3,4)
	(1,3)	R →	(3,3)
	(1,2)	(2,2)	(3,2)
	G ←	(2,1)	(3,1)
		(2,0)	

Lowest-Cost First Search

- Consider parent cost!

```
n = state(init)
frontier.append(n)
while(frontier not empty)
  n = pull state from frontier
  visited.append(n)
  if n = goal, return solution
  for all actions in n
    n' = a(n)
    if n' not visited
      priority = heuristic(goal,n')
      frontier.append(priority)
```

What cost heuristic could we add?

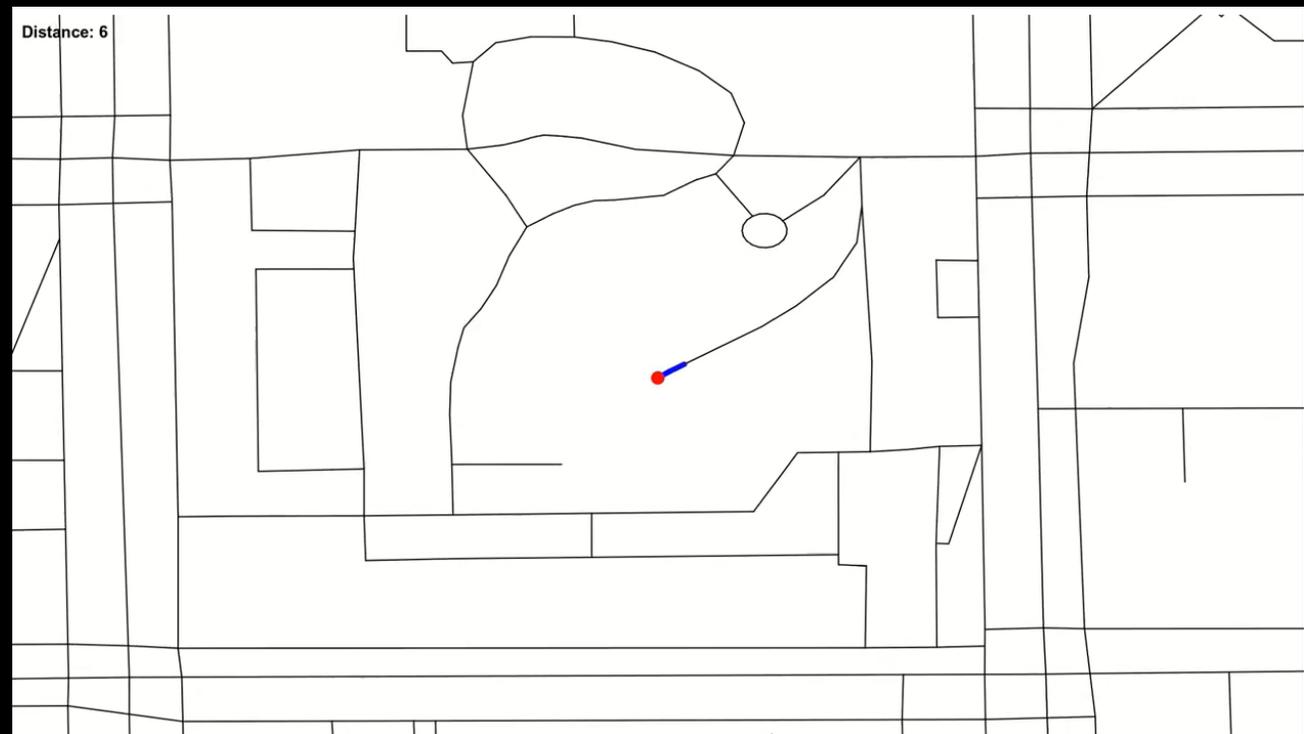
- Go straight, cost 1
- Turn one quadrant, cost 1

	(1,4)	(2,4)	(3,4)
	(1,3)	R →	(3,3)
	(1,2)	(2,2)	(3,2)
		↓	
	G ←	(2,1)	(3,1)
		(2,0)	

Lowest-Cost First Search

- Is it *complete*?
 - Yes, as long as path costs are positive
- What is the time complexity?
 - $O(b^m)$
- What is the space complexity?
 - $O(b^m)$

<https://www.youtube.com/watch?v=t7UjtzqIXSA>

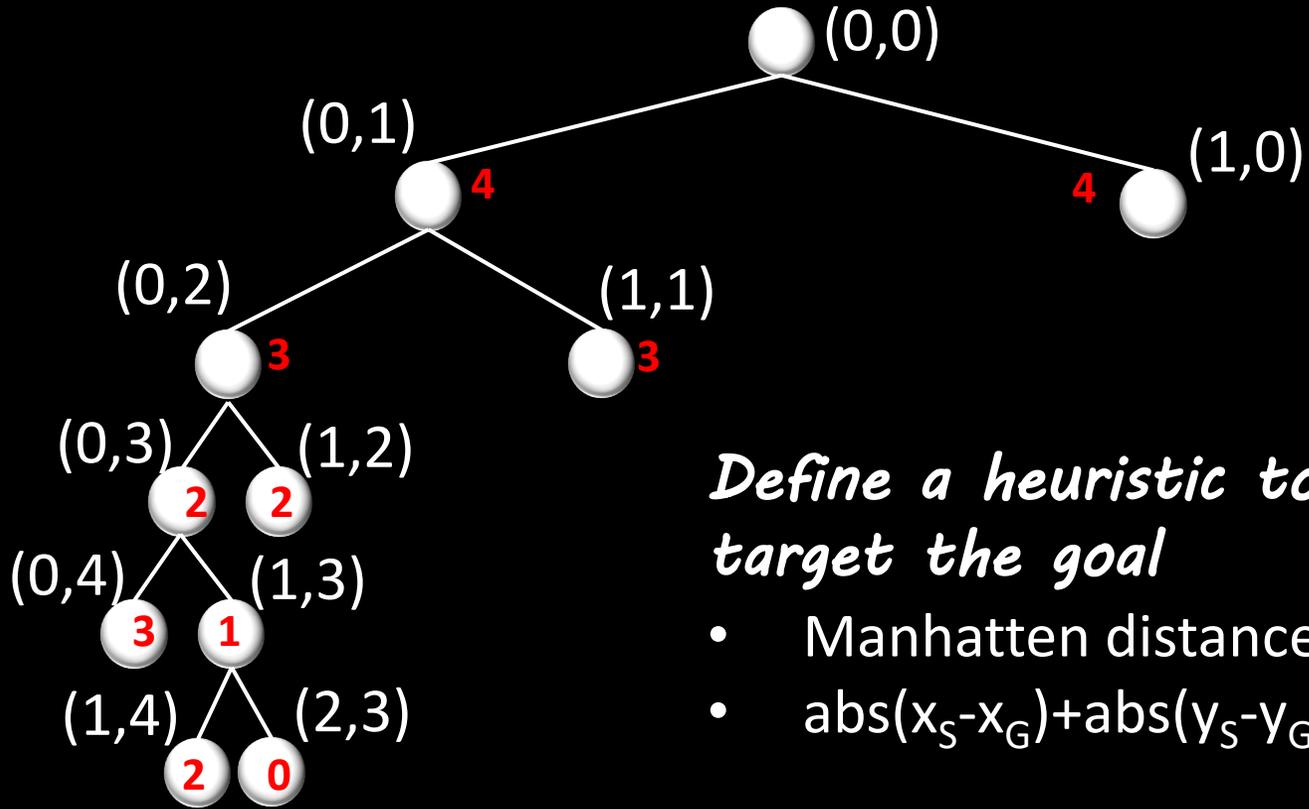


Could we be smarter?

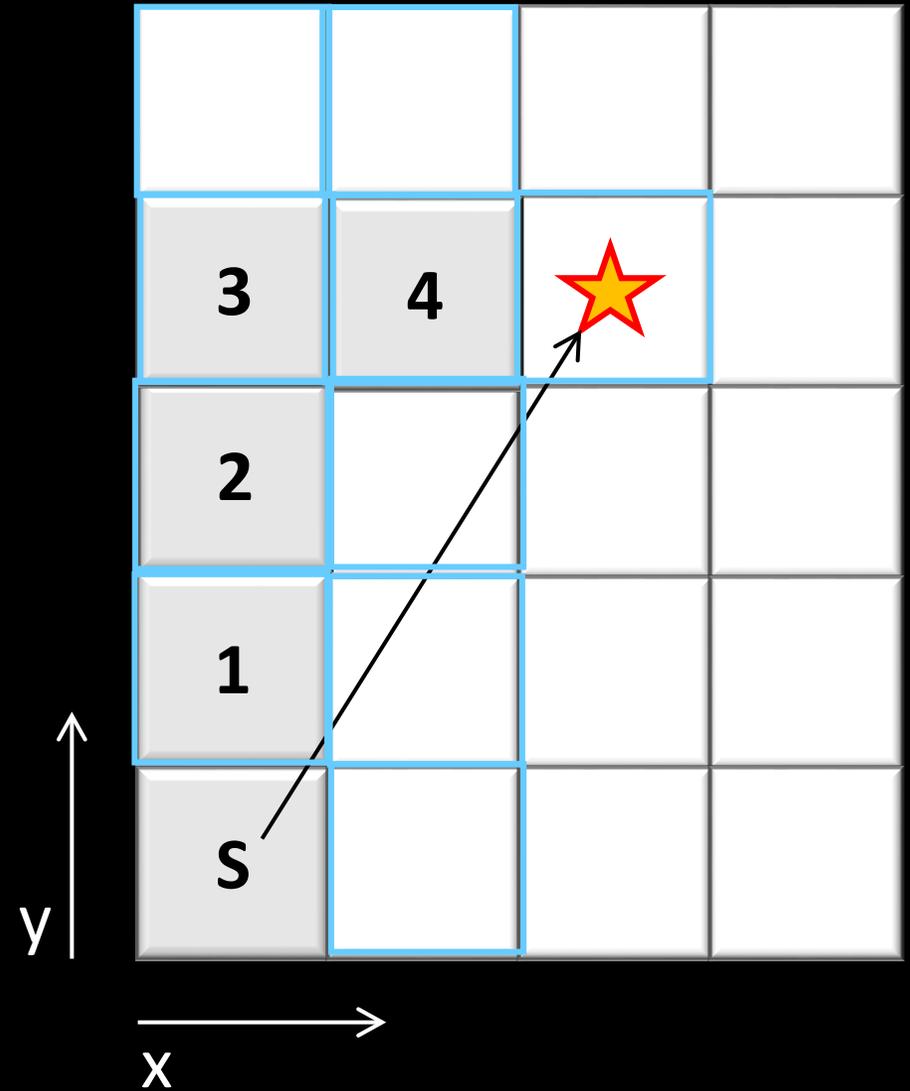
- *Sure, you know the graph and you know the goal is!*
- ...Informed search
 - Consider parent cost, and..
 - ..estimate the shortest path to the “goal”
- Assign a value to the frontier
 - Pick frontier closest to the goal (minimize distance)

Informed Search

- Greedy Search



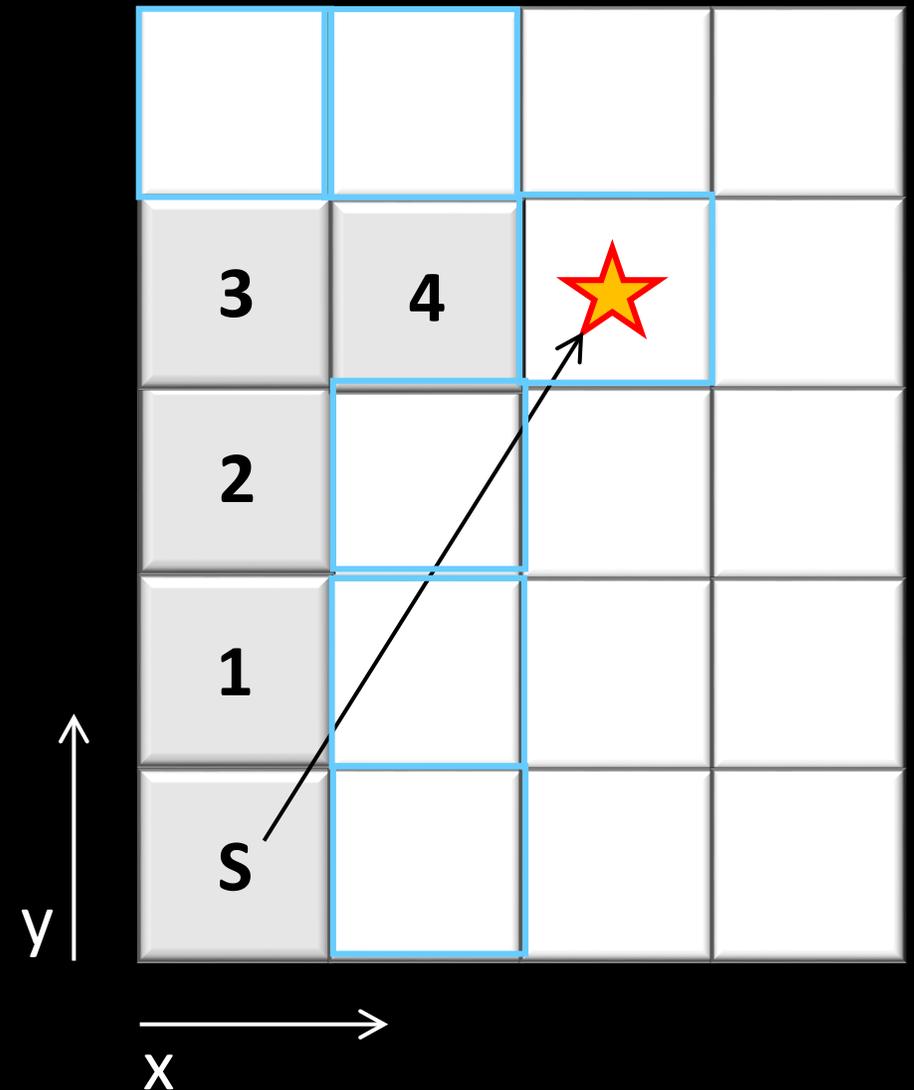
Search order: N, E, S, W



Informed Search

- Greedy Search

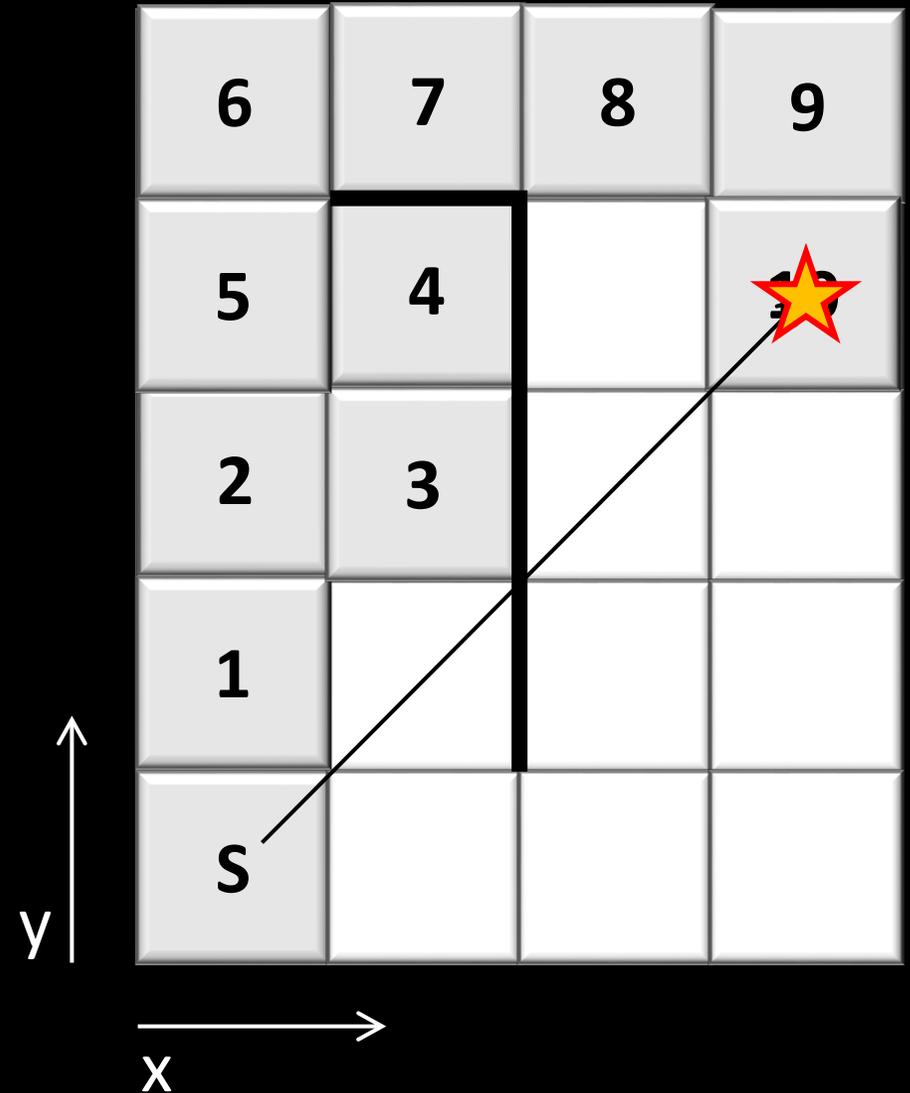
```
n = state(init)
frontier.append(n)
while(frontier not empty)
    n = pull state from frontier
    visited.append(n)
    if n = goal, return solution
    for all actions in n
        n' = a(n)
        if n' not visited
            priority = heuristic(goal,n')
            frontier.append(priority)
```



Informed Search

- Greedy Search
 - Complete?
 - No
 - Time complexity?
 - $O(b^m)$
 - Space complexity?
 - $O(b^m)$
 - Optimal?
 - no...

Search order: N, E, S, W



Search Algorithms, General

- Breadth First Search
 - *Complete* and optimal
 - ...but searches *everything*
- Lowest-Cost First Algorithm *Considers parent cost*
 - *Complete and optimal*
 - ...but it wastes time exploring in directions that aren't promising
- Greedy Search *Considers goal*
 - *Complete (in most cases)*
 - ...only explores promising directions

Can we do better?

*A**

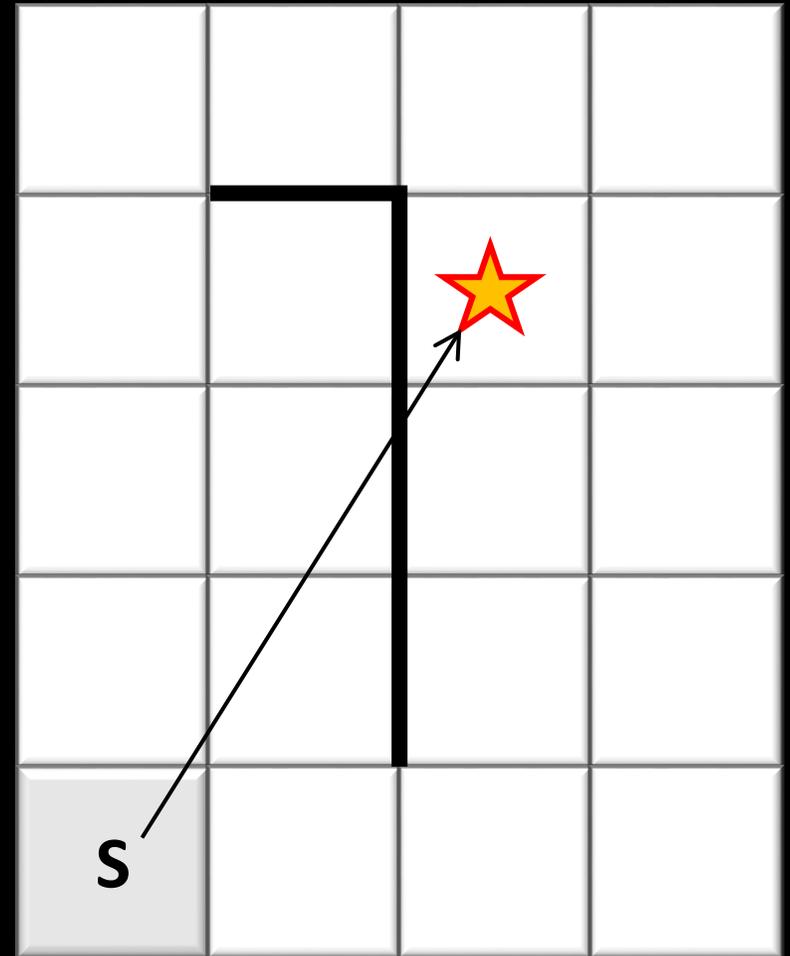
Informed Search

- A* (“A-star”)

```
n = state(init)
frontier.append(n)
while(frontier not empty)
  n = pull state from frontier
  if n = goal, return solution
  for all actions in n
    n' = a(n)
    if ((n' not visited or
        (visited and n'.cost < n_old.cost))
        priority = heuristic(goal,n') + cost
        frontier.append(priority)
        visited.append(n')
```

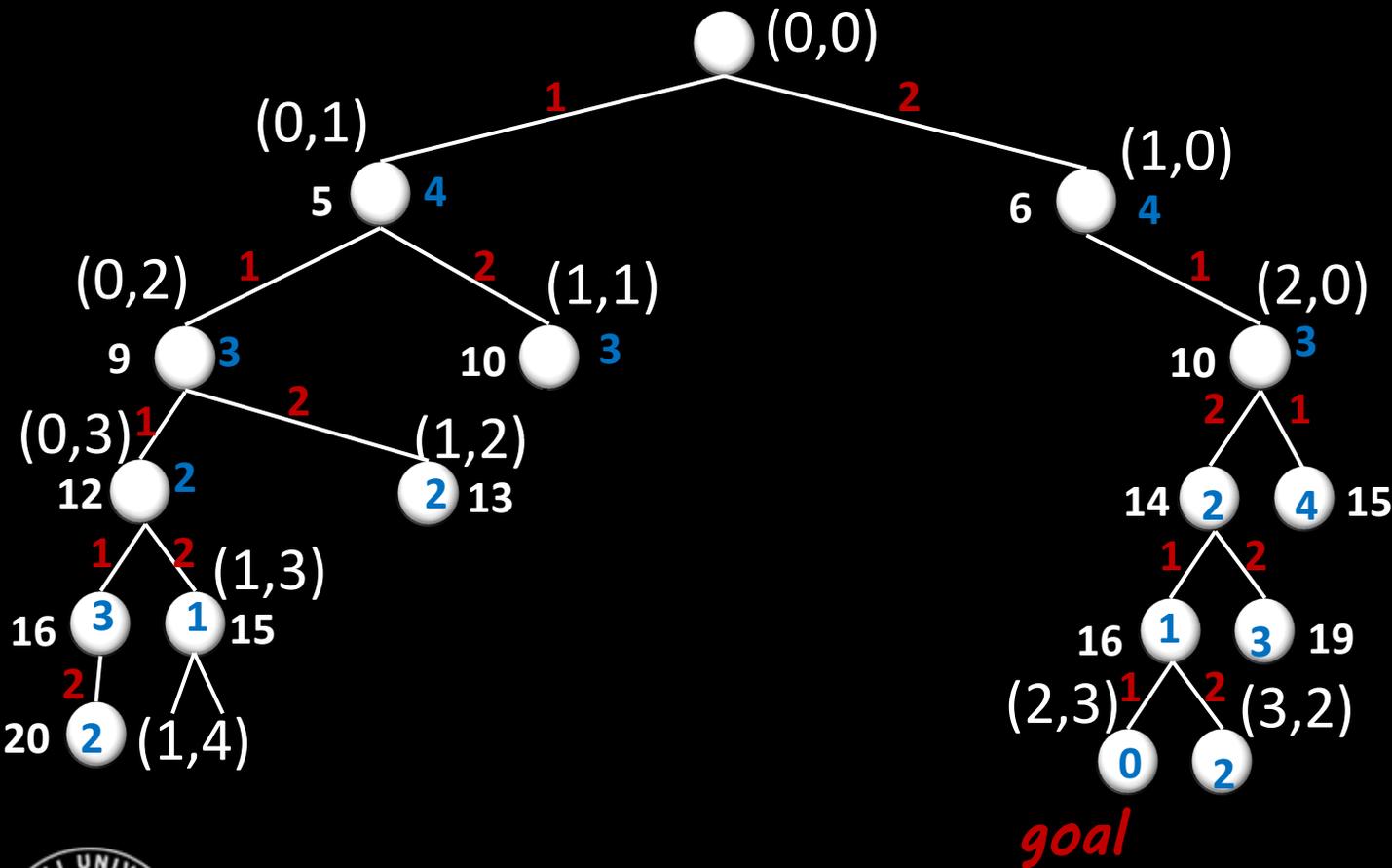
Search order: N, E, S, W

Find a treasure



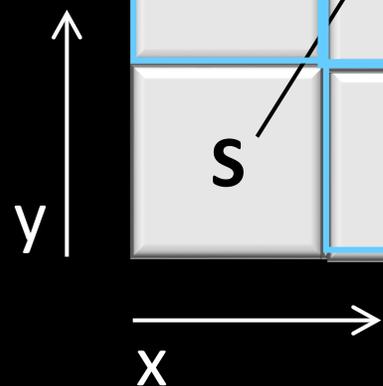
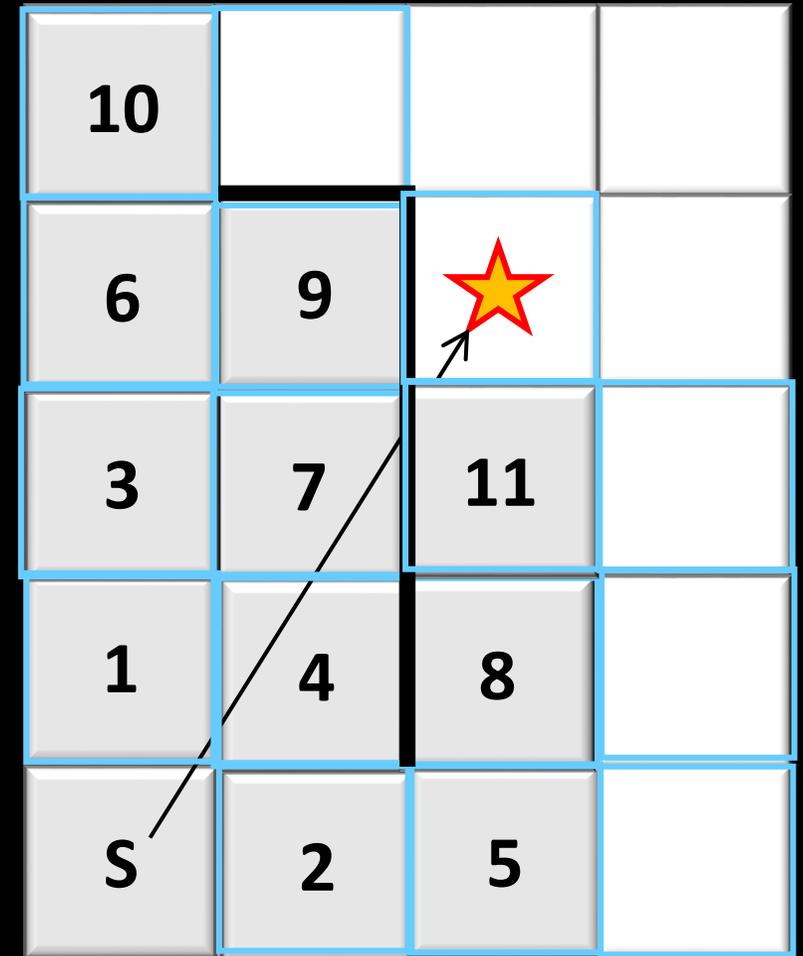
Informed Search

- A* ("A-star")
 - Cost and goal heuristic



Search order: N, E, S, W

Find a goal



A* Search

- What if the heuristic is too optimistic?
 - Estimated cost $<$ true cost
- What if the heuristic is too pessimistic?
 - Estimated cost $>$ true cost
 - No longer guaranteed to be optimal
- What if the heuristic is just right?
 - Pre-compute the cost between all nodes
 - Feasible for you?

admissible heuristic

inadmissible heuristic



Summary

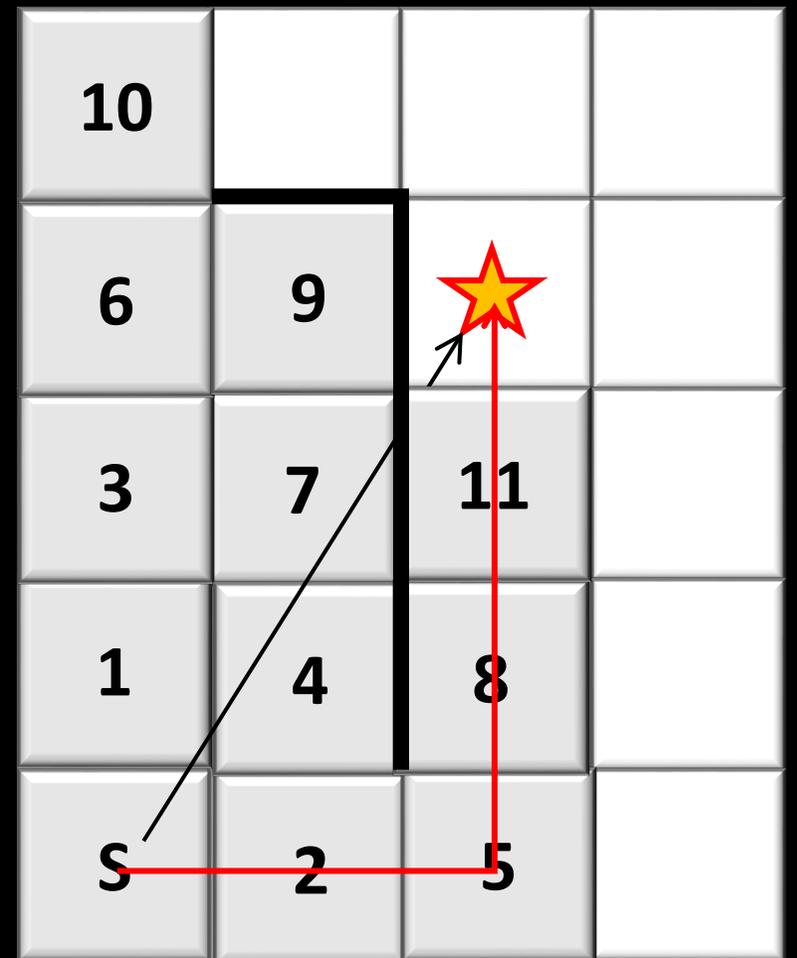
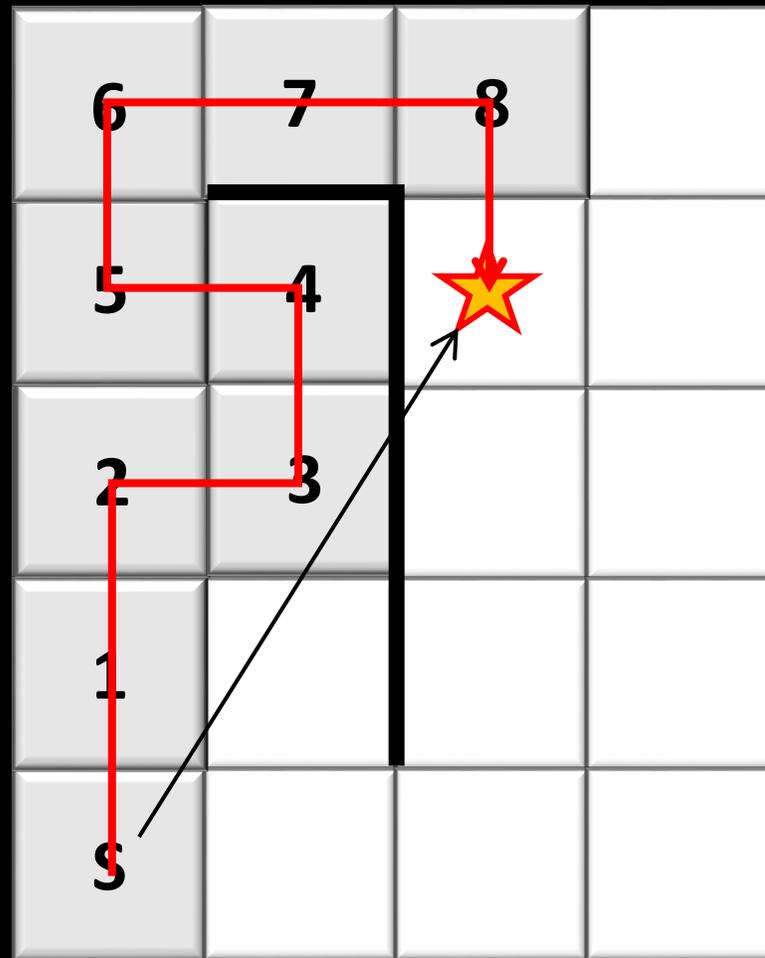
LCFS

minimum path

Greedy

A*

*minimum path
and efficient*



Icons for navigation modes: Car, Bus, Walking, Bicycling, Airplane.

Upson Hall, 124 Hoy Rd, Ithaca, NY 14850

Cornell Dairy Bar, 411 Tower Rd, Ithaca, NY 14850

Add destination

OPTIONS

- Send directions to your phone
- via Hoy Rd
18 min
0.8 mile
DETAILS
- via Hoy Rd and NY-366 E/Dryden Rd
21 min
1.0 mile
- via Hoy Rd and Tower Rd
22 min
1.0 mile

